

Kap 23: Formelle språk og grammatikker

$$A = \{ a, b \}$$

$$L = \{ a, aa, aab, \dots \}$$

• union : $L = \{ a, b \}$ $L \cup \Pi = \{ a, b, bb, c \}$

$$\Pi = \{ \underbrace{bb}_a, c \}$$

ikke kommutativ

• konkatenering : $L = \{ \underline{a}, \underline{b} \}$
 $\Pi = \{ \underline{bb}, c \}$

$$L\Pi = \{ abb, ac, bbb, bc \}$$

$$\Pi L = \{ bba, bbb, ca, cb \}$$

↓
 identitetelement : $\{ \Lambda \}$

$$\frac{L\Pi \neq \Pi L}{L^3 = LLL = \{ aa, ab, ba, bb \} \{ a, b \} =}$$

• tillutning : $L = \{ a, b \}$

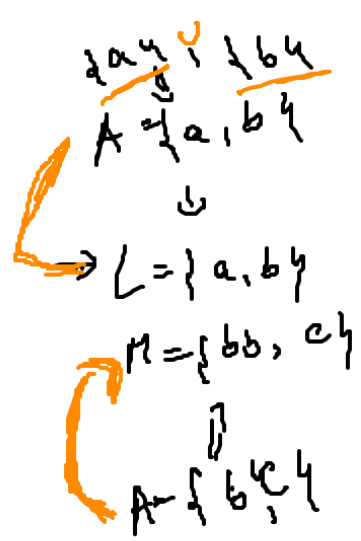
$$\underline{L} \Rightarrow L^0 \cup L^1 \cup L^2 \cup \dots \cup L^u$$

$\{ \Lambda \}$ $\{ a, b \}$

$$\Lambda a = a \Lambda = a$$

$$L^2 = LL = \{ a, b \} \{ a, b \} = \{ aa, ab, ba, bb \}$$

$\emptyset \Rightarrow \{\}$
 $L\emptyset = \emptyset L = \emptyset$
 $\Rightarrow \{ \wedge \} \Rightarrow \{ \wedge \} \Rightarrow L^0$
 \uparrow
 $S \wedge = \wedge S = \underline{S}$



$L^0 \cup L^1$
 \uparrow

Def: Reg. språk $\rightarrow A = \{a, b\}$

$\bullet) \emptyset, \{ \wedge \}, \{ a^k \}$ for hvert $a \in A$

\Rightarrow Semantikk $\bullet) \text{ hvis } L \text{ og } M \text{ er reg. språk, så er } \underline{LM}, L \cup M, L^*, M^*$

Def: Reg. uttrykk

$\bullet) \emptyset, \wedge, a$ for hvert $a \in A$

$\bullet) \text{ hvis } L \text{ og } M \text{ er reg. uttrykk så er } L, LM, L^*, M^* \rightarrow (L)$

\Downarrow
 $\text{tolke}(LM) = \text{tolke}(L) \text{ tolke}(M)$
 $\text{tolke}(LM) \Rightarrow \text{tolke}(L) \cup \text{tolke}(M)$
 Støtaker



$$\underline{01^*} = \{0, 1, 11, 111, \dots\}$$

$(011)^*$ ⇒ alle strings over alphabet $\{a, b\}$

$$0^*1^* = \{1, 0, 00, 000, \dots, 1, 11, 111, \dots\}$$

$$\rightarrow \underline{(011)^*} = \{1, 0, 1, \underbrace{00, 01, 10, 11}_{2 \text{ tegn}}, \dots\}$$

1 tegn

$$\{0, 1, 2\}$$

$$(0112)^*$$

union $\rightarrow (01)^* = \{1, 01, 0101, 010101, \dots\}$

konkat

\downarrow

$(011)^*$ same som $\underline{0^*1^*}$ ⇒ $\{1, 0, 1, 11, 111, \dots\}$

like to

$$L = \{ \lambda, a, ab, abb, abbb, \dots \}$$

$$\lambda | a b^* \Rightarrow \cancel{ab^*}$$



$$ab^*$$

$$a \lambda, b, abb, \dots$$

$$\cup$$

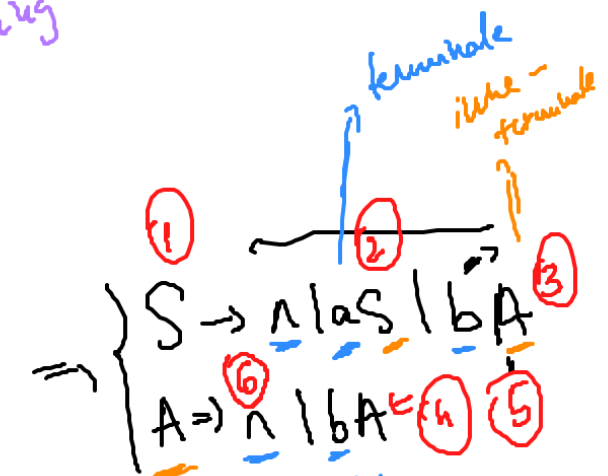
$$a$$

$$S \rightarrow \lambda$$

$$S \rightarrow aS | bA$$

$$A \rightarrow \lambda | bA$$

$$\underbrace{a, ab, abb}_{aS | b}$$



1. read
 S \Rightarrow aS \Rightarrow abA \Rightarrow abbbA \Rightarrow abbb

$L = \{ \overset{\uparrow}{\Lambda}, ab, aabb, aaabbb, \dots \} \Rightarrow$ regulär
 \uparrow
 $a^n b^n$

$\mathcal{U}_1 S \Rightarrow \Lambda \mid a \overset{\downarrow}{S} b \mid ab$
 \uparrow
 $aabb$

$\mathcal{U}_2 S \Rightarrow \Lambda \mid a \overset{\downarrow}{S} b$
 \uparrow
 $aabb$

$S \Rightarrow aSb \Rightarrow aabb$
 \uparrow
 \uparrow

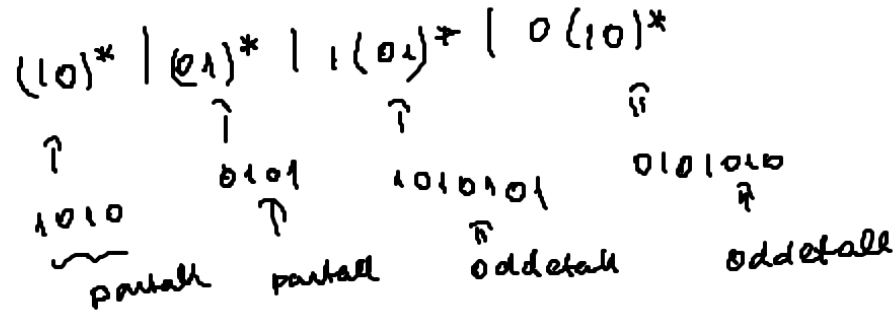
$a^* (ab)^* b^*$

$a^* b^*$
 \Downarrow
 $\Rightarrow a^* b^* \quad (ab)^*$
 \downarrow
 abb, aab, \dots

$ab, abab, ababab$
 \uparrow
 $(ab)^*$

23.11 b) {0, 1}

01, 010, 10, 101



a) 0 eller 11

$(0(111))^* (0(1))^*$

↑
with some strings for alphabet