Reading, Analyzing, and Presenting Scientific Papers

IN5420/9420, SPRING 2023

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Classifying the type of contribution

- A tutorial or a textbook
- Elegant formalization of a problem that captures reality
- New or improved solution to an important and interesting problem
- New generic technique for analysis/solution/evaluation
- New analysis
- Better evaluation
- Survey

2

Important questions to consider

- Target application domain
  - What makes the work so suitable for this type of applications?
  - How general or easy is it to extend?
- Is the solution practical?
  - Example: the solution requires system engineer to define an exponential number of...
- Is the evaluation comprehensive and adequate?
  - Example: dependable system in failure-free settings or adaptive system with most parameters fixed

3

Important questions to consider (cont´d)

- What is the motivation for this work?
  - The authors certainly claim that it is important. What makes you believe it?
- Does the contribution and results correspond to the motivation? Do all pieces fit together?
- Where did the work fall short of your wild imagination? What extension would you like to see the most?
- Asking questions is always an option
  - Teachers, supervisors, colleagues
  - Even if you do a presentation alone

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Guidelines for reading papers

- Read the paper in three passes
  - Pass I: the abstract, section titles, techniques, and references
  - Pass II: read the paper but prioritize, skip the proofs and elements less important for understanding the contribution
  - Pass III: try to re-implement the paper in your mind
    - Identify hidden assumptions
    - Identify weak and strong points
    - Missing comparison with related work
    - Potential issues with the experimental or analytical techniques
  - Refer to "How to read a paper" https://www.albany.edu/spatial/WebsiteFiles/ResearchAdvices/how-to-read-a-paper.pdf

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Guidelines for student presentations

- Present authors' claim regarding the work
  - Problem studied
  - Contributions of the paper
  - Relation to the state of the art
  - Depending on type of contribution and research method
    - Definitions, Requirements, Design, Algorithms, Implementations, Experiments, Evaluation, etc
  - Main results
  - Conclusion
  - Refer to http://www.stanford.edu/~jacksonm/present.pdf

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Guidelines for student presentations (cont´d)

- Criticism
  - Asking and concluding on the important questions
  - Guidelines: “The Task of the Referee” by Alan J Smith, UCB
  - Guidelines: Section 4 in “Writing Reviews for Systems Conferences” by T. Roscoe
- Questions for discussion!
- Summary and conclusion of criticism

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For presentations you will be better off skipping

- Formal parts
  - Unless it is a work on new proof techniques or formal methods
  - Formal notation and definitions
    - important for written papers
    - replace with intuition and examples
  - Theorem proofs and mathematical derivations
    - focus on the result and intuition behind the proof instead

8

For presentations you will be better off skipping (cont´d)

- Lengthy lists of related work
  - Instead present categories
  - Or one representative for each
- Fine-grain implementation details
  - Present the central part or implementation idea
    - Skip optimizations
  - Provide overview of pseudo-codes
    - and if require, present a few central lines
- Overly detailed architectural schemes
  - Papers typically include such schemes with dozens of entities and arrows, and annotations in a tiny font

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Guidelines for preparing questions

Avoid following questions:

- Questions with answer well-known and easily found in the literature
  - E.g., the average time between blocks in Bitcoin
- Questions absolutely unrelated to technology
  - E.g., velocity of a given cryptocurrency
  - It is ok, however, to talk about the use and popularity of technology because it gives motivation for the technology
- Questions that we can have no hope to reason about
  - E.g., will cryptocurrencies be in use in 100 years

Good questions:

- Understanding of nontrivial points in the topic
- Technological outlook for the near future
- Comparison between studied technologies

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Guidelines for writing reviews

- Hypothesis and Context
  - brief problem description
  - contribution
  - summarize the central ideas, research methodology and supporting points
- Novelty of the research
  - a new perspective on the topic,
  - applicable to real-world situations,
  - comparison with other approaches
  - providing concrete examples of where the approach succeeds or fails.
- Assessment
  - your opinion of the research

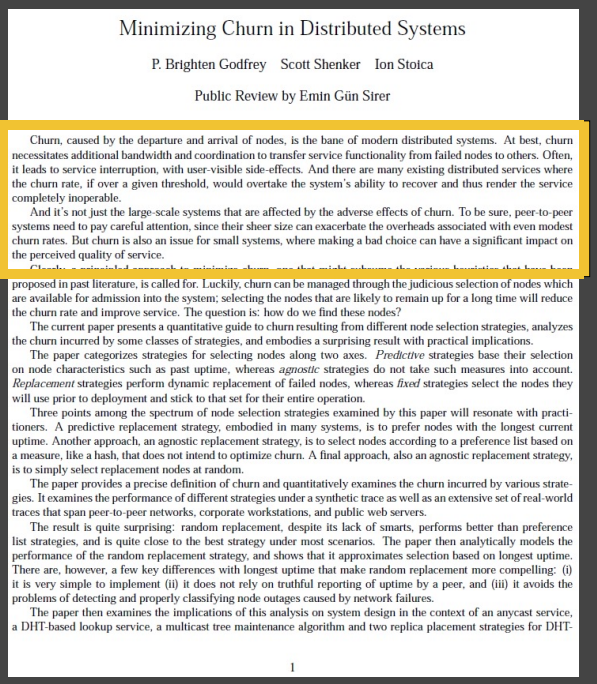
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Walk-through a sample review

- Research Paper: "Minimizing Churn in Distributed Systems"
- Public review available at Sigcomm 2006

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Motivation/Importance of the Problem



Minimizing Churn in Distributed Systems

P. Brighten Godfrey Scott Shenker Ion Stoica

Public Review by Emin Gün Sirer

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Clearly, a principled approach to minimize churn, one that might subsume the various heuristics that have been proposed in past literature, is called for. Luckily, churn can be managed through the judicious selection of nodes which are available for admission into the system; selecting the nodes that are likely to remain up for a long time will reduce the churn rate and improve service. The question is: how do we find these nodes?

The current paper presents a quantitative guide to churn resulting from different node selection strategies, analyzes the churn incurred by some classes of strategies, and embodies a surprising result with practical implications.

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The paper provides a precise definition of churn and quantitatively examines the churn incurred by various strategies. It examines the performance of different strategies under a synthetic trace as well as an extensive set of real-world traces that span peer-to-peer networks, corporate workstations, and public web servers.

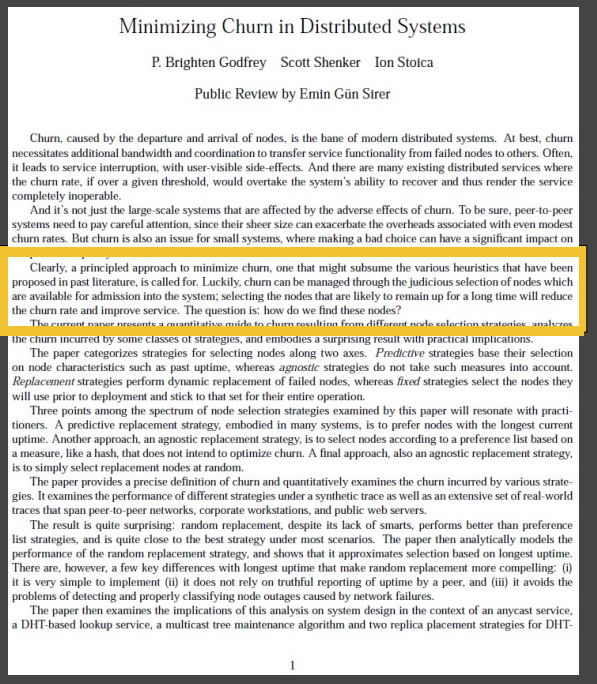
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Definition of the Problem



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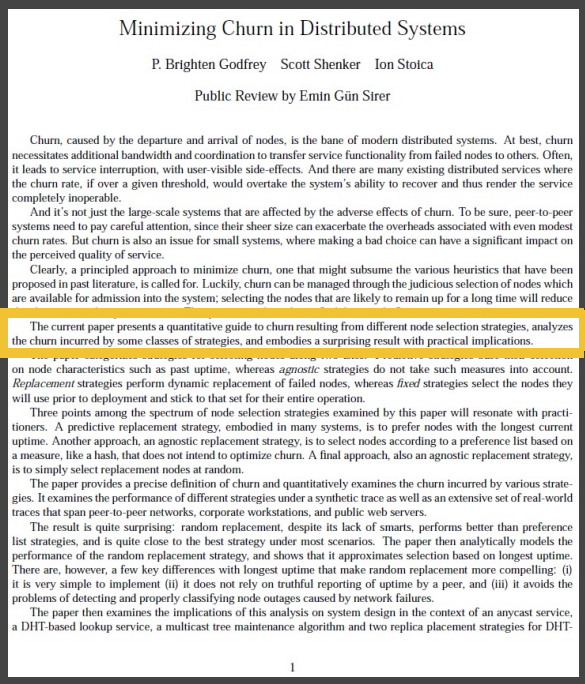
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Overview of Contributions



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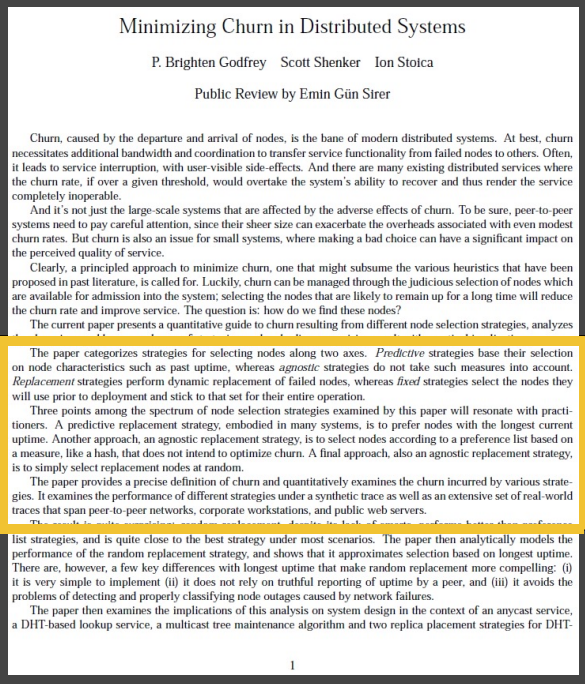
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Research Methodology



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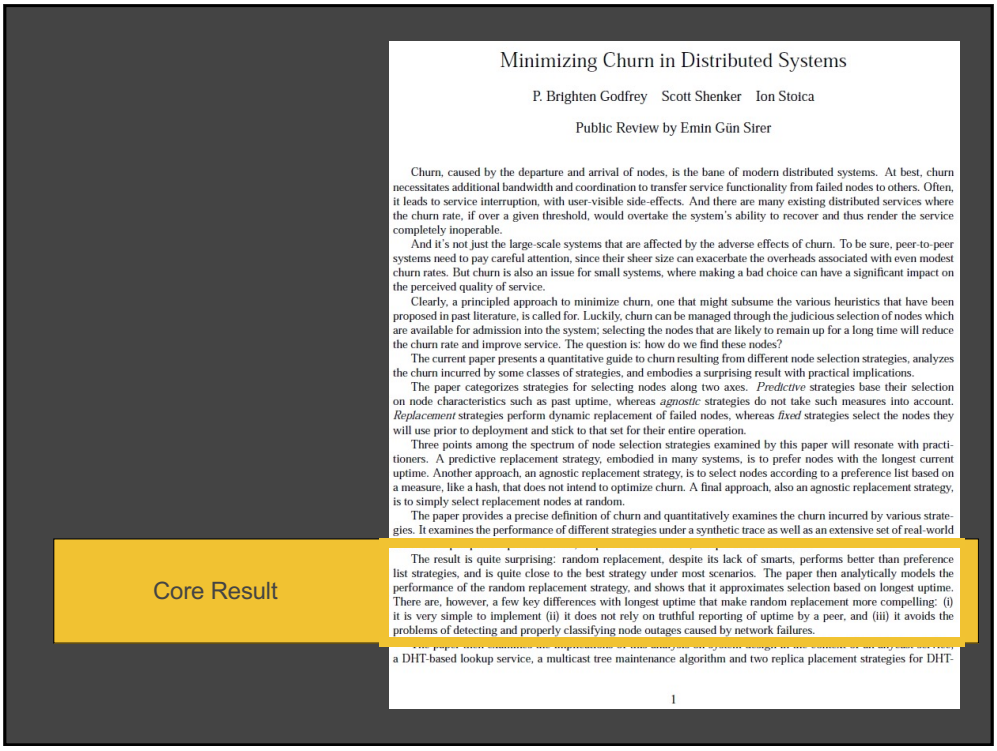
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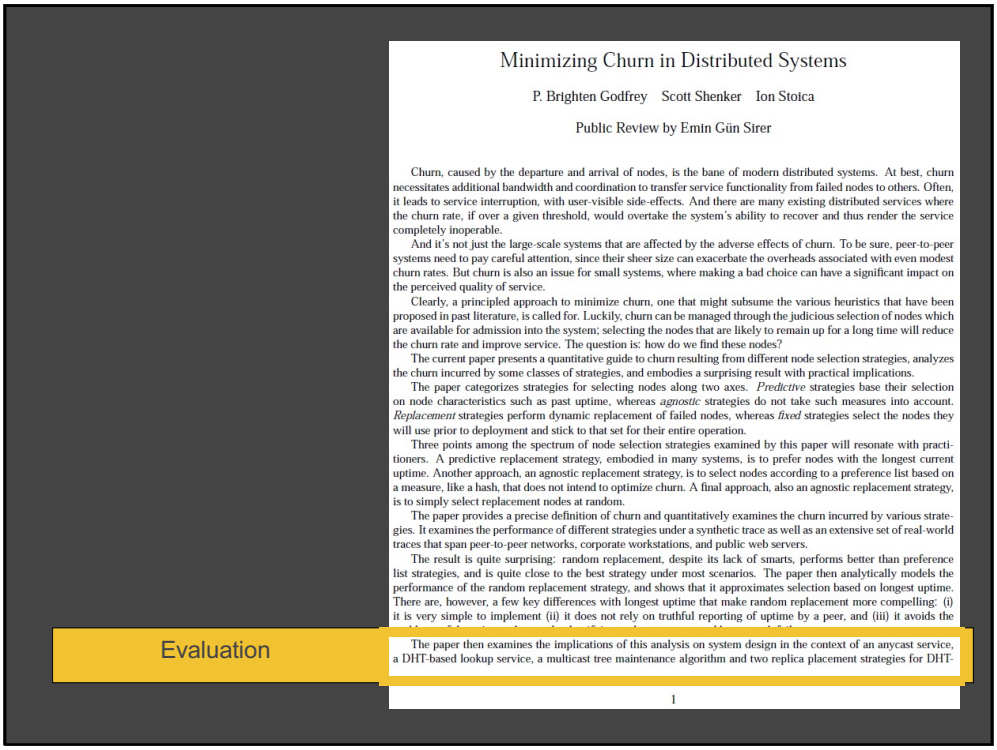
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Core Result

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Evaluation

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Evaluation

based storage services. The last scenario extends the quantitative analysis to a setting where nodes carry long-term state. For each scenario, the paper examines end-to-end performance metrics, and discusses how randomization can be used to transform commonly used strategies that, in essence, equate to a preference list into node selection criteria that reduces churn by approximating random selection. The changes required are straightforward and small, yet have an appreciable impact on the performance of the systems.

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It is rarely the case that system designers can be told to do less to improve performance. This paper shows that less is, in this context, better.

2

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Open Issues/Problems

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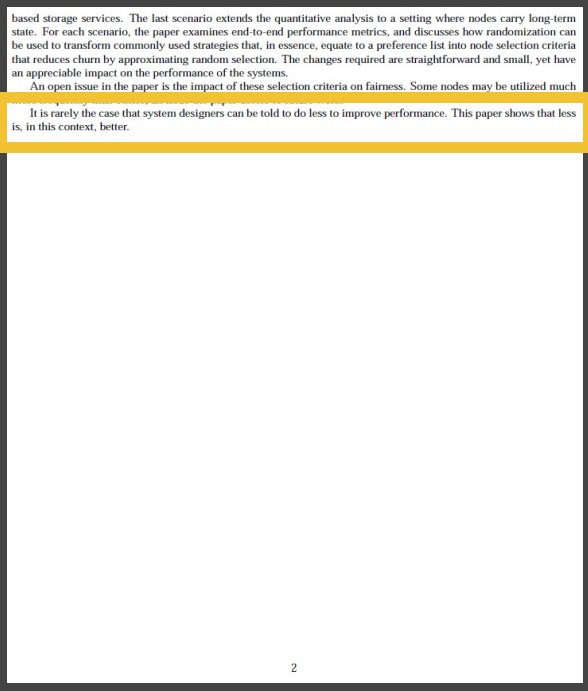
An open issue in the paper is the impact of these selection criteria on fairness. Some nodes may be utilized much more frequently than others; an issue the paper defers to future work.

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Take-away Message



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Reading material

- A general collection of links on https://www.cs.columbia.edu/~hgs/etc/writing-style.html
- Reading and Presenting Research Papers, http://www.cs.rpi.edu/academics/courses/spring99/robotics/paperdiss.html
- Review
  - ACM Computing Reviews guidelines, http://www.computingreviews.com/Reviewer/RVguidelines2015.pdf
  - "The Task of the Referee" by Alan Jay Smith, especially the "Evaluating a research paper" section
- Presentation tips
  - http://www.stanford.edu/~jacksonm/present.pdf
  - http://ksuweb.kennesaw.edu/faculty/rbrow211/papers/presenting-research.html
  - Winston, P. H., "How to speak"

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