

Introduction to Data Communication

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(based on slides from Kjell Åge Bringsrud and Carsten Griwodz)

Introduction

Goal

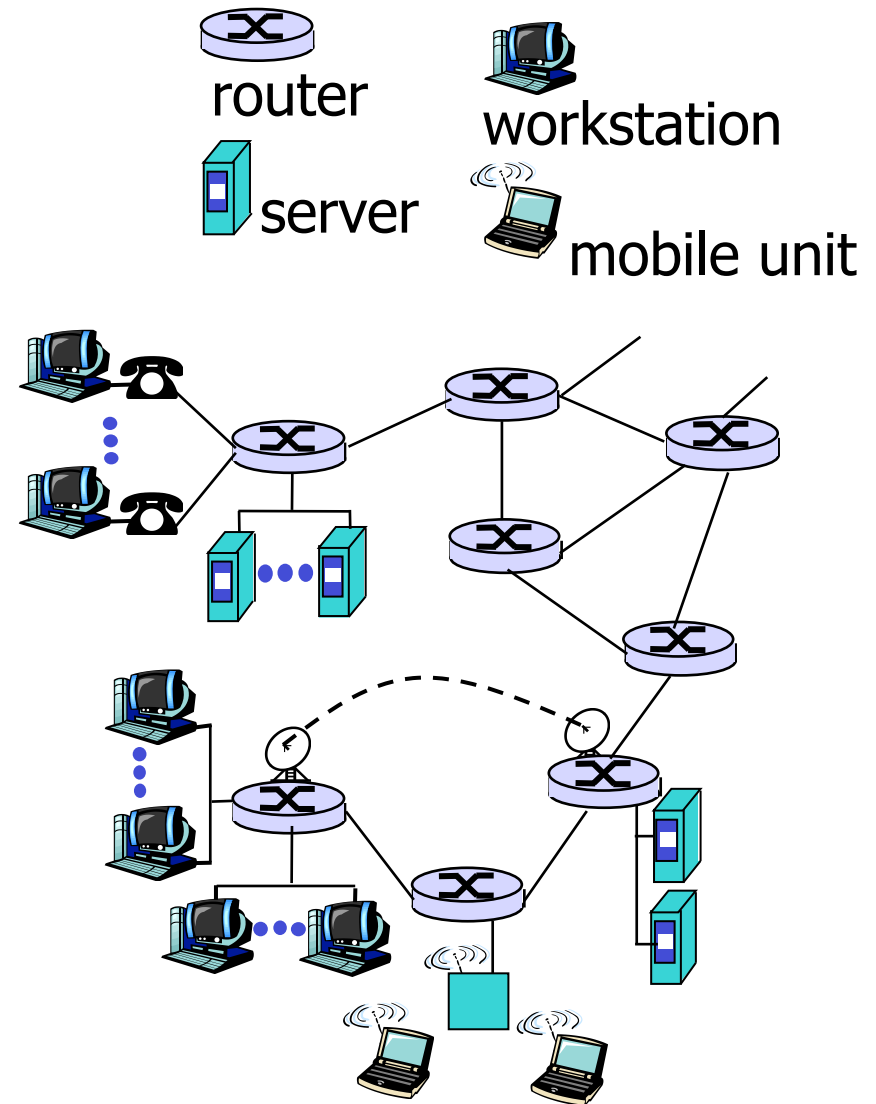
- ❑ Give an overview of the topic
- ❑ Approach
 - Descriptive
 - Use Internet as example

Content

- ❑ What is the Internet?
- ❑ What is a protocol?
- ❑ End systems
- ❑ Access network and physical media
- ❑ Core networks
- ❑ Throughput, delay, and loss
- ❑ Protocol layers, service models
- ❑ Backbones, NAP' er, ISP' er
- ❑ History

What is the Internet?

- ❑ Millions of interconnected devices: **host computers, end systems**
 - PCs, workstations, servers
 - PDAs, telephones, fridges ... which run
- network applications**
- ❑ *Communication links*
 - Fiber, copper, radio, satellite
- ❑ **Routers**
 - passing on packets of data through the network



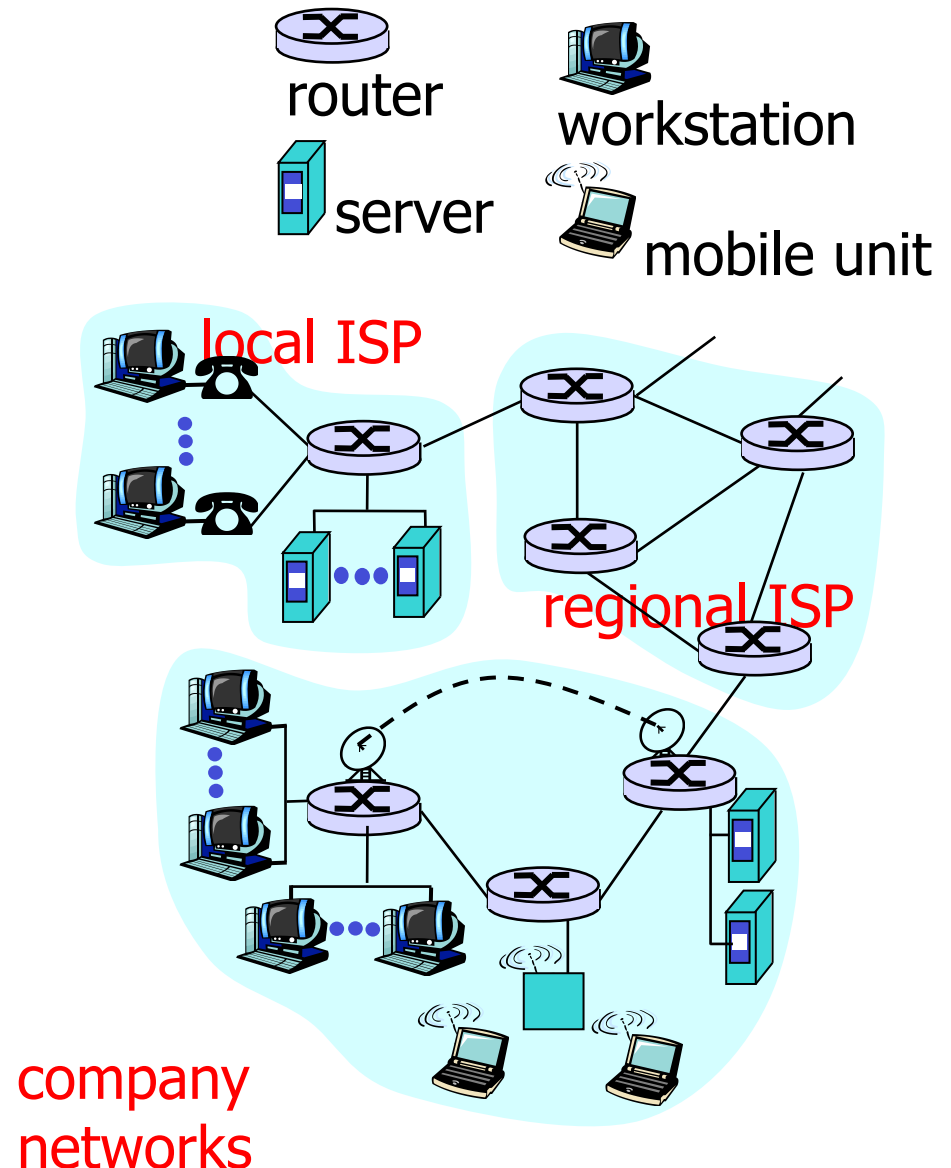
What is the Internet?

❑ *Internet: “network of networks”*

- Partly hierarchical
- ISPs: Internet Service Providers
- Public Internet versus private intranet

❑ *Protocols*

- Control sending, receiving of messages
- E.g., TCP, IP, HTTP, FTP, PPP



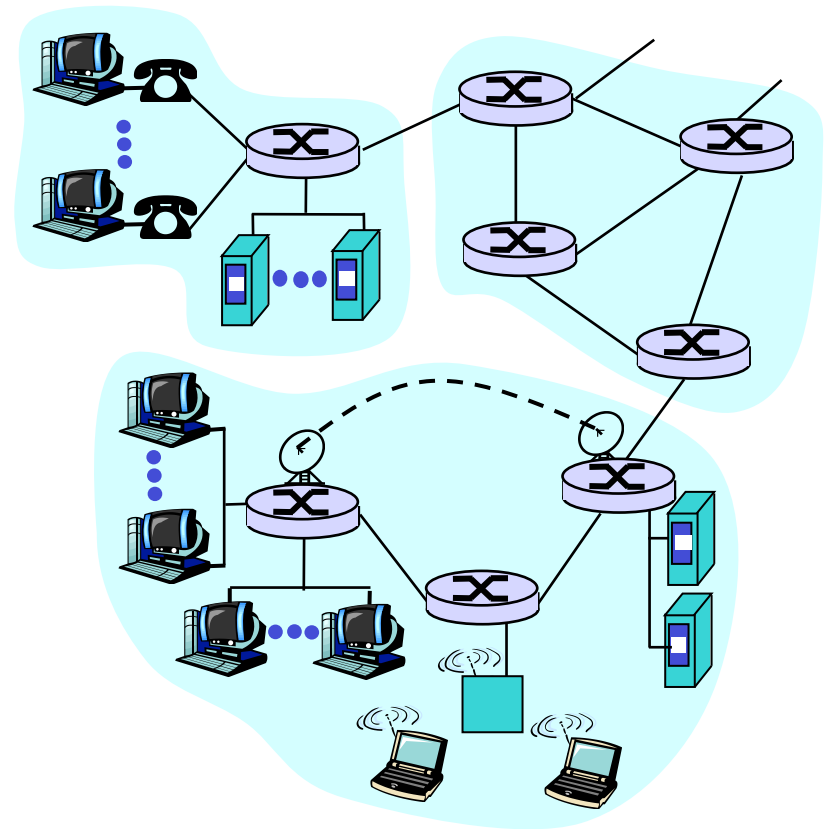
What is the Internet from a service view?

❑ Communication *infrastructure*

- Allows distributed applications:
- WWW, email, games, e-commerce, database, elections,
- More?

❑ Internet standards:

- RFC: Request for comments, e.g. TCP is RFC 793
- IETF: Internet Engineering Task Force



End systems

❑ End systems

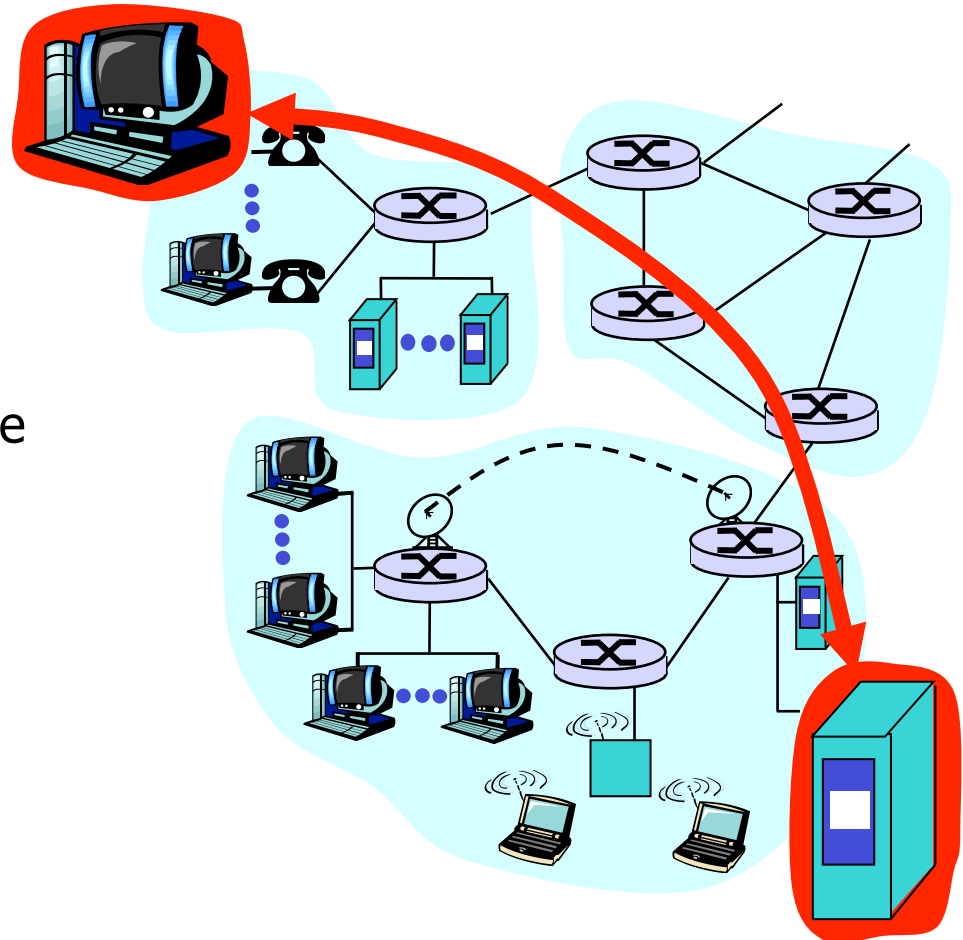
- Run application programs
- E.g., web browser, web server, email
- At “the edge” of the net

❑ Client/server model

- Clients ask for, and get a service from the servers
- E.g. WWW client (browser)/server; email client/server

❑ Peer-to-peer model

- Interactions are symmetrical
- E.g. telephone conferences



What is a protocol?

Human protocols:

- ❑ “What time is it?”
- ❑ “I have a question”
- ❑ Formal phrases...

... are special
“messages” that are
sent, which lead to ...
... defined events or
actions when the
message is received

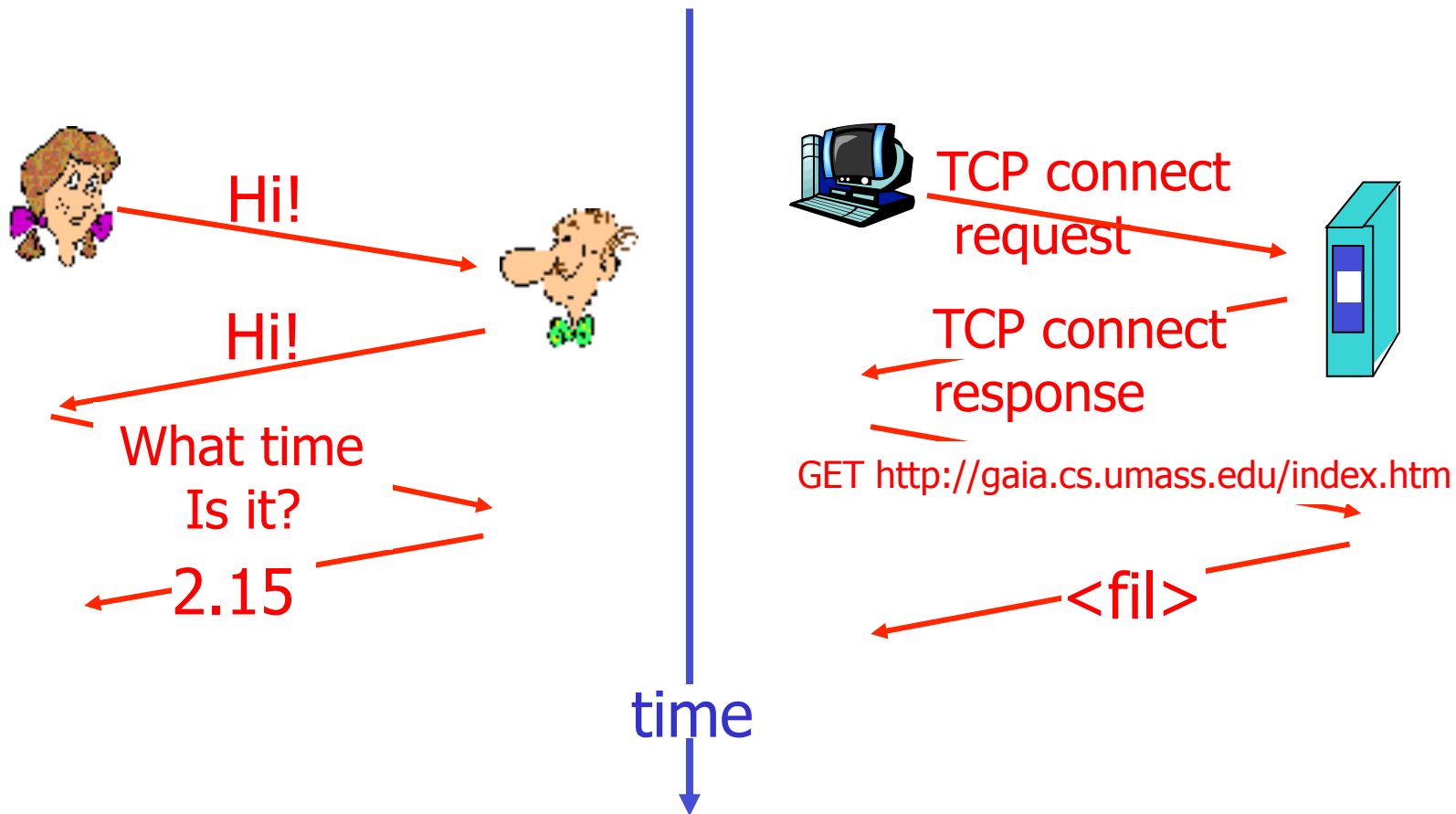
Network protocols:

- ❑ Machine instead of people
- ❑ All communication activity in the Internet is controlled by protocols

Protocols define formats, order of sending and receiving of messages, and the actions that the reception initiates.

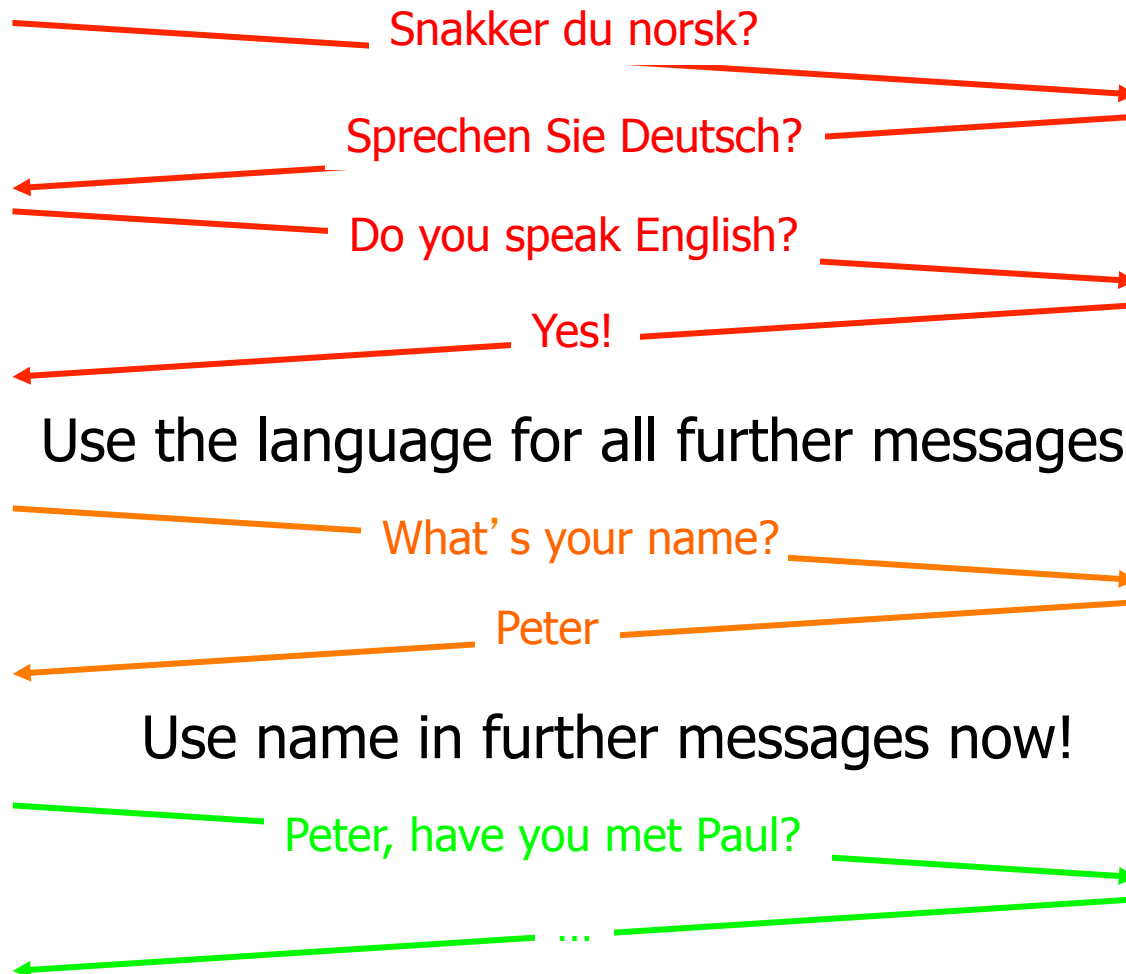
What is a protocol?

A human protocol and a computer protocol:



What are protocol layers?

Several layers of communication



time

What are protocol layers?

Networks are complex

- ❑ Many parts:
 - Hardware, software
 - End systems, routers
 - Links of different kinds
 - Protocols
 - Applications

Question:

Is it possible to *organize* the structure of a network?

Or at least our discussion of networks?

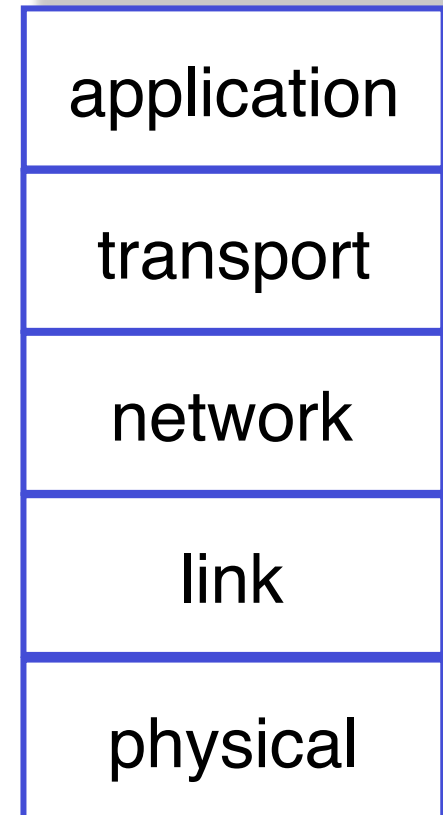
Why layering?

Management of complex systems:

- ❑ Modularisation simplifies
 - Design
 - Maintenance
 - Updating of a system
- ❑ Explicit structure allows
 - Identification of the individual parts
 - Relations among them
- ❑ Clear structure: layering
 - Layered **reference model**
 - Goal: different implementation of one layer fit with all implementations of other layers

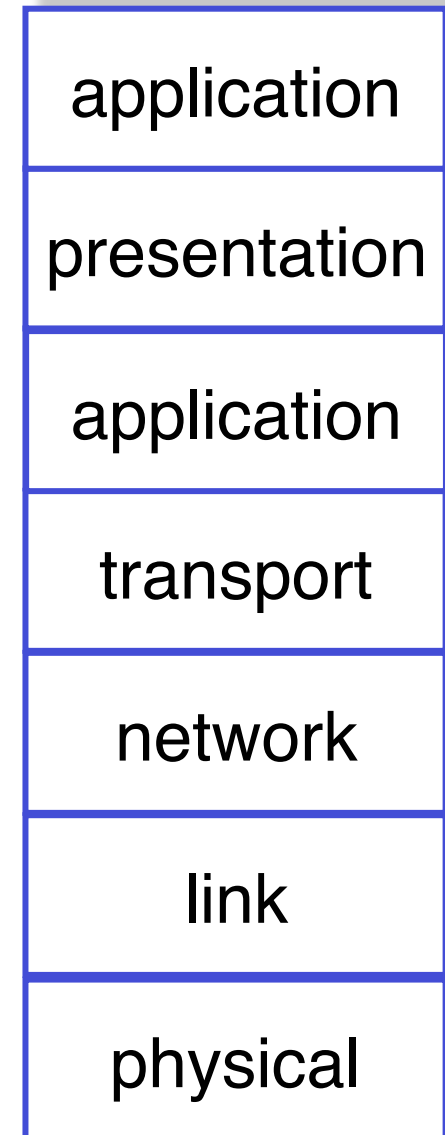
TCP/IP - protocol stack

- ❑ **application:** supports network applications
 - ftp, smtp, http
 - Your applications
- ❑ **transport:** data transfer from end system to end system
 - TCP, UDP
- ❑ **network:** finding the way through the network from machine to machine
 - IP
- ❑ (data) **link:** data transfer between two neighbors in the network
 - ppp (point-to-point protocol), Ethernet
- ❑ **physical:** bits “on the wire”



OSI - model

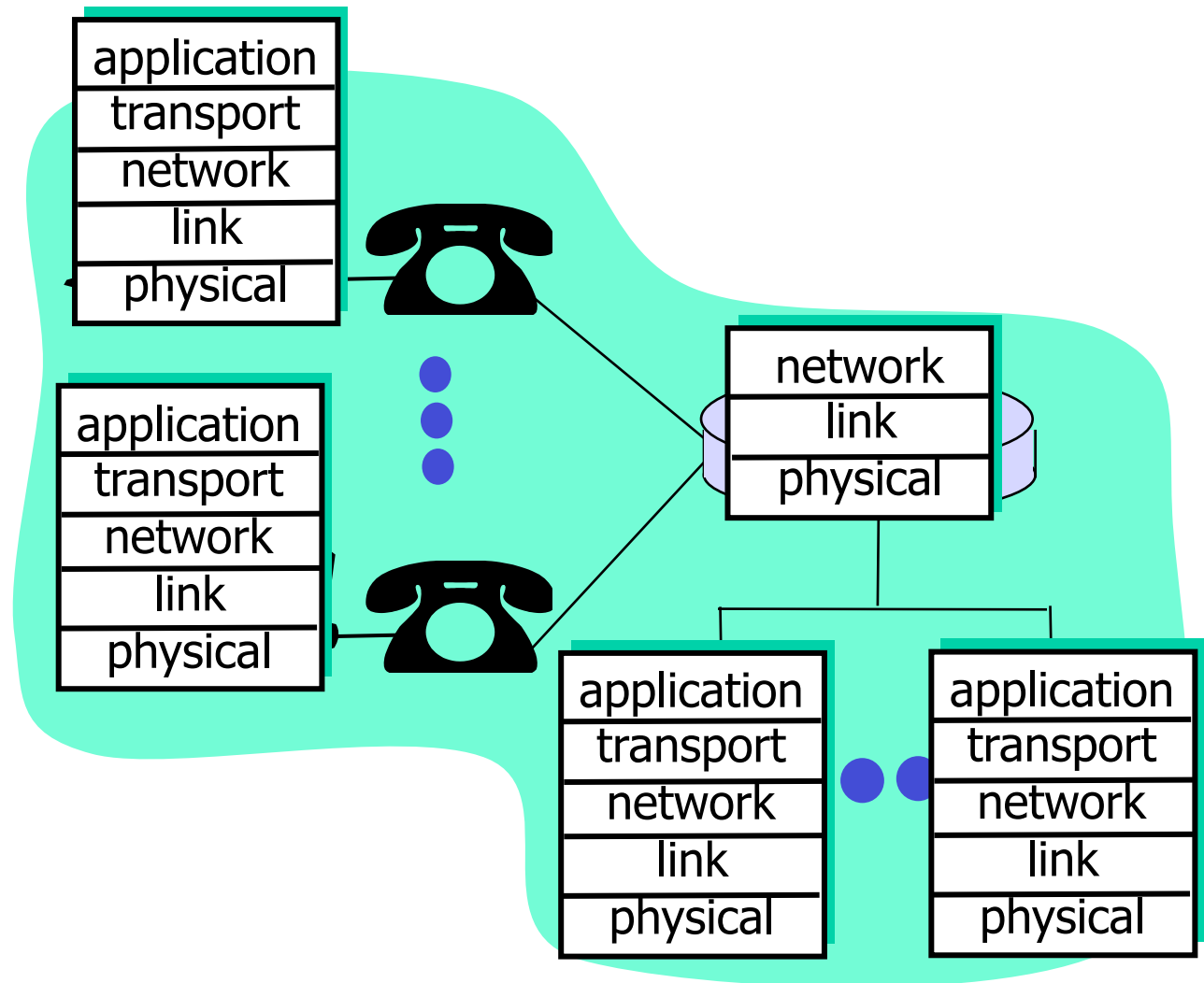
- ❑ A standard for layering of communication protocols
 - Open Systems Interconnection
 - by the ISO – International Standardization Institute
- ❑ Two additional layers to those of the Internet stack
- ❑ **presentation:** translates between different formats
 - XML, XDR
 - provides platform independence
- ❑ **session:** manages connection, control and disconnection of communication sessions
 - RTP



Layering: logical communication

Each layer:

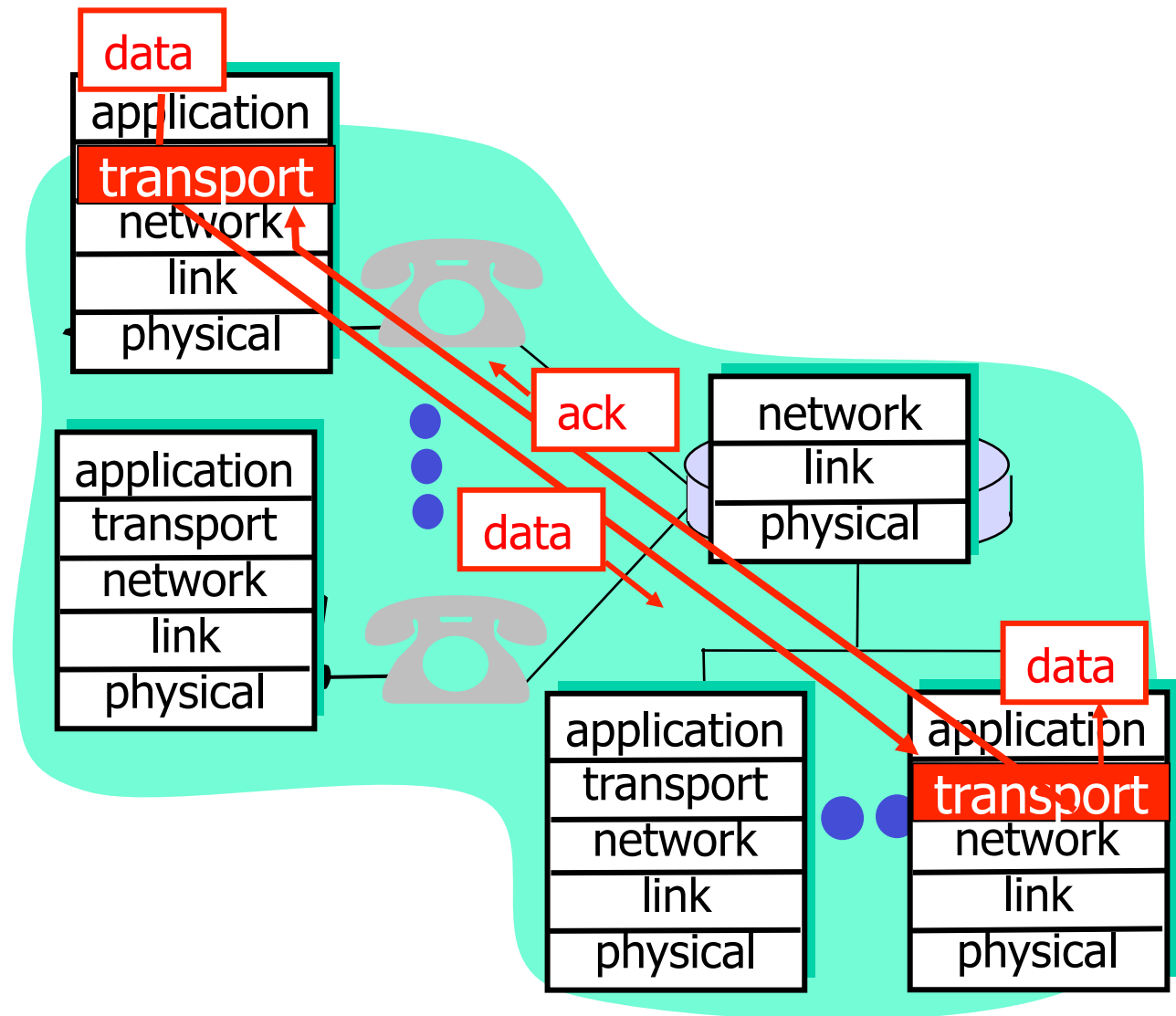
- ❑ distributed
- ❑ “units” implement functionality of each layer in each node
- ❑ Units execute operations, and exchange messages with other units of the same layer



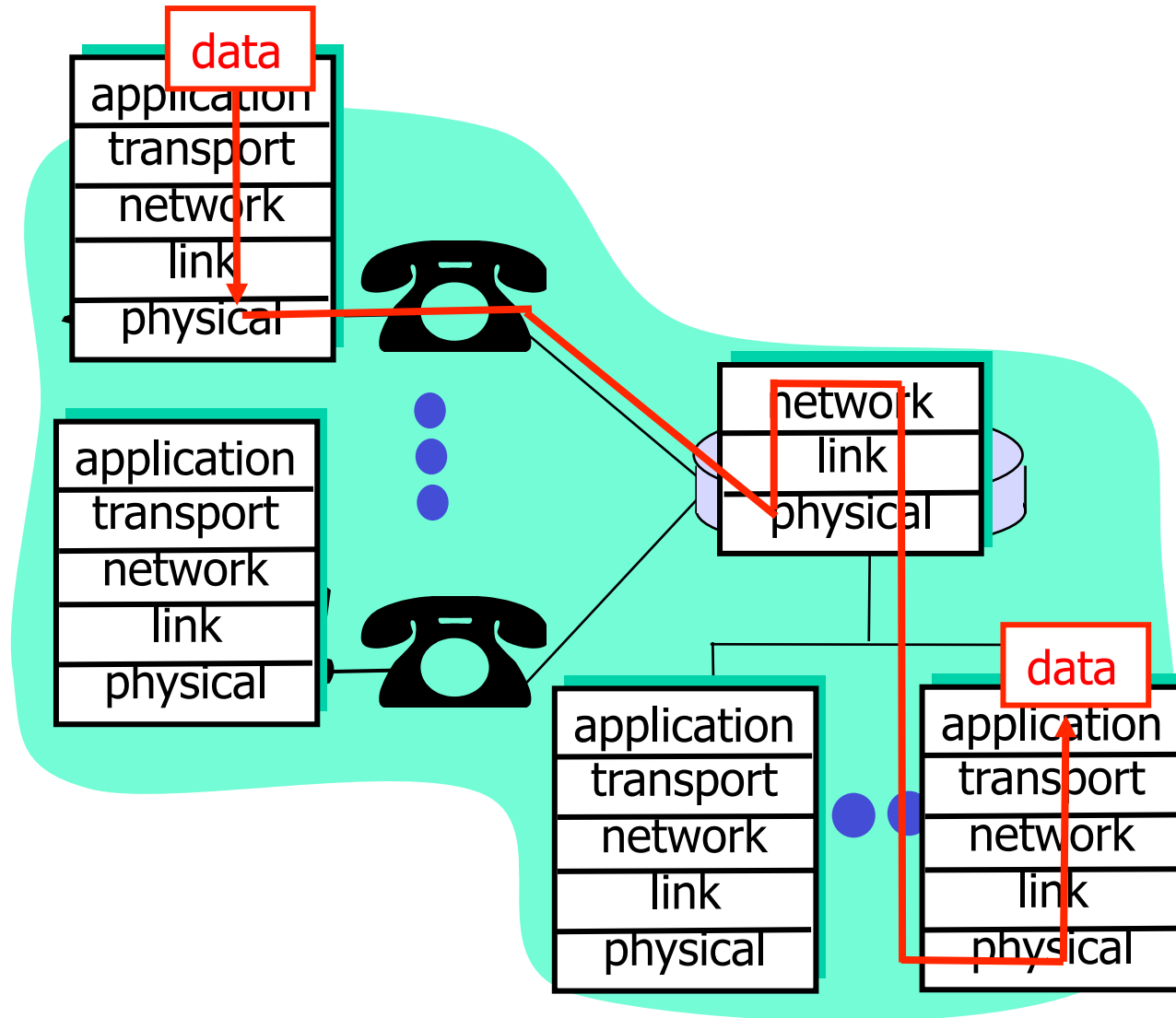
Layering: *logical* communication

E.g. transport

- ❑ Receive data from the application
- ❑ Add receiver address, reliability check, information to create a “datagram”
- ❑ Send datagram to the transport layer in the receiver node
- ❑ Wait for “ack” from the transport layer in the receiver node
- ❑ Analogy: post office



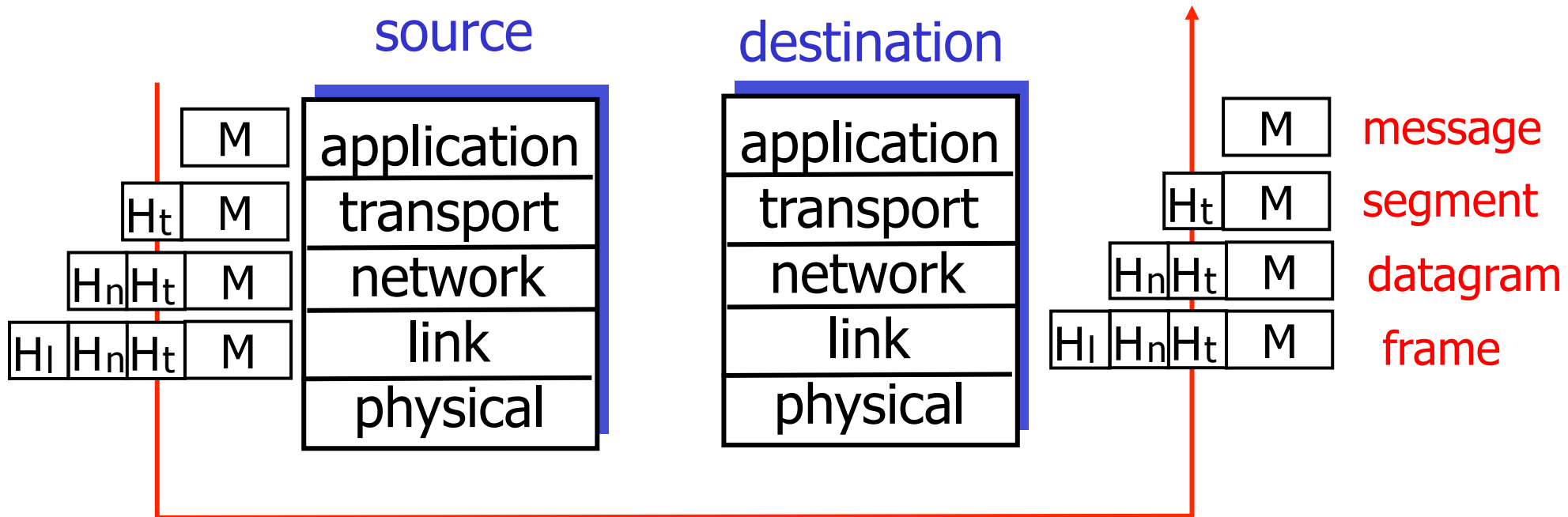
Layering: physical communication



Protocol layer and data

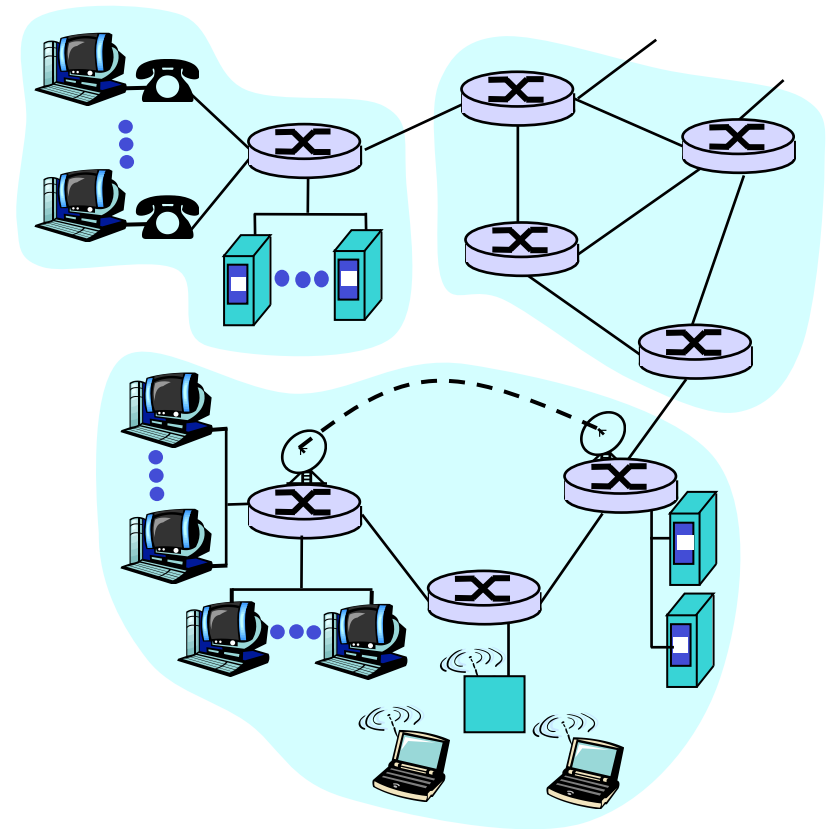
Each layer takes data from next higher layer

- ❑ Adds header information to create a new data unit (message, segment, frame, packet ...)
- ❑ Send the new data unit to next lower layer



A closer look at network structures

- ❑ **End systems**
 - applications and host computers
- ❑ **Access network, physical medium**
 - Communication links
- ❑ **Core networks**
 - Routers
 - Network of networks



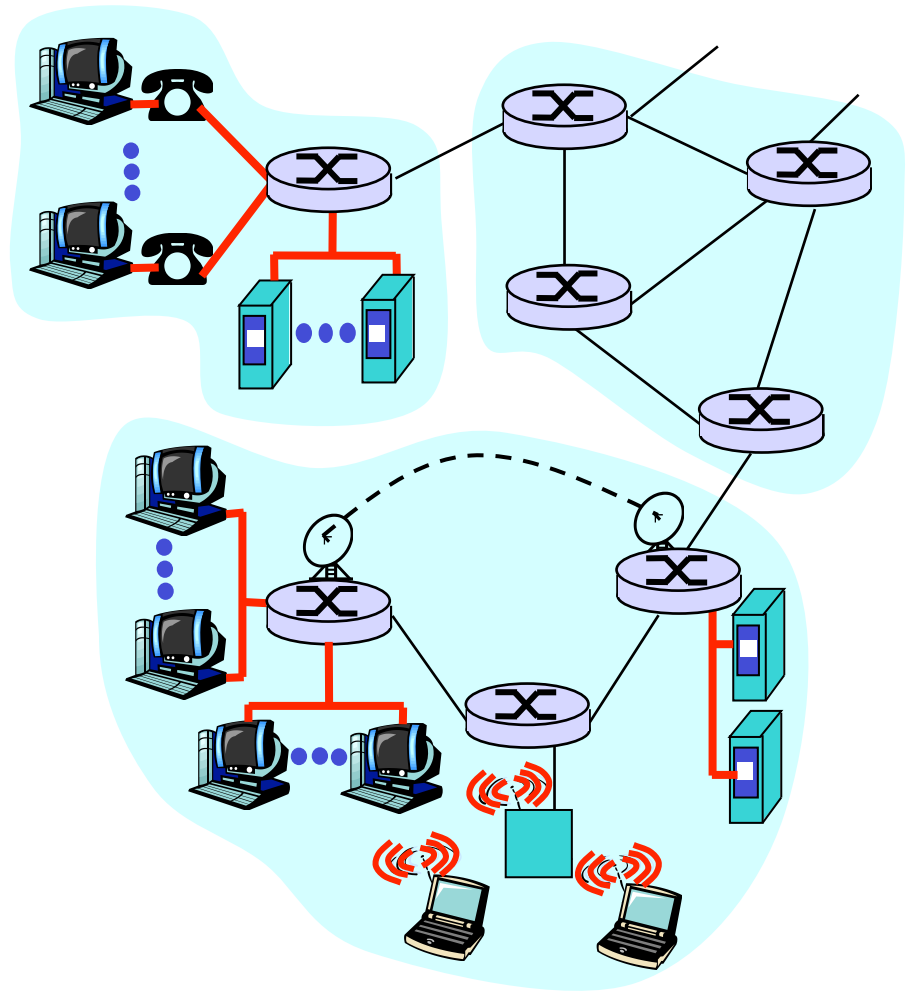
Access network and physical media

How to connect end systems to edge routers?

- ❑ Home network
- ❑ Company network (schools, companies)
- ❑ Mobile access network

Keep in mind when choosing a technology:

- ❑ Bandwidth?
- ❑ Shared or dedicated medium?



Home network: point to point

❑ Dial-up via modem

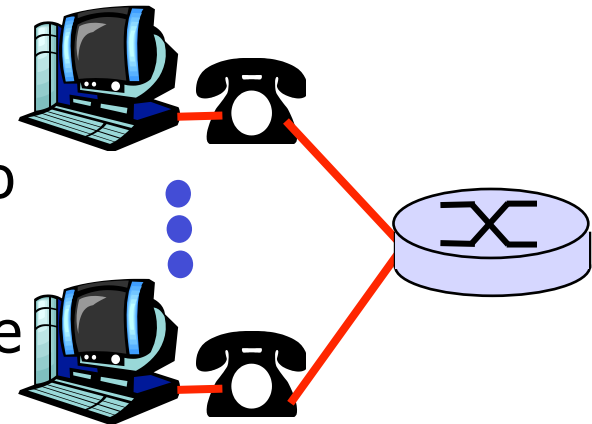
- Up to 56Kbps direct access to the router (at least in theory)

❑ ISDN: integrated services digital network

- 128Kbps purely digital connection to the router

❑ ADSL: asymmetric digital subscriber line

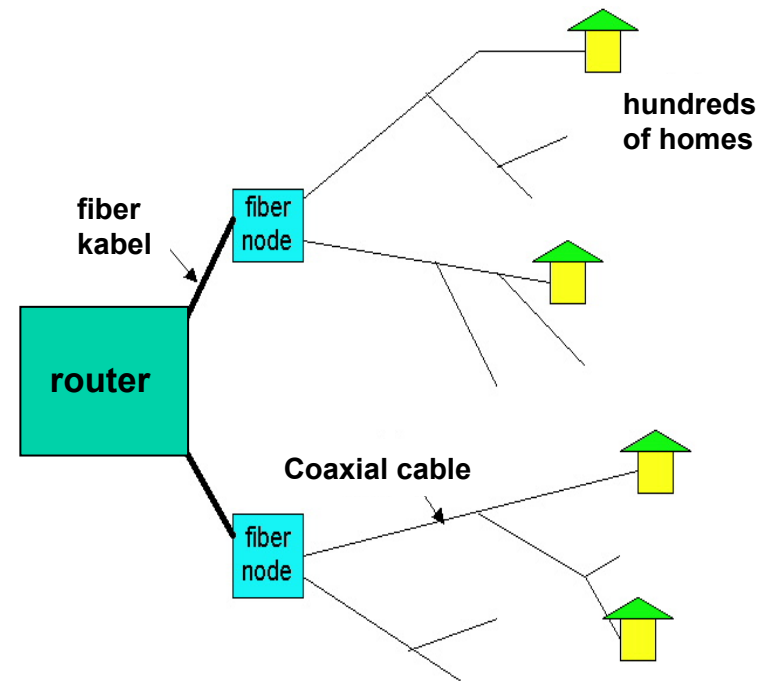
- Up to 5 Mbps *uplink* (home-to-router, ver. ADSL2++, in development)
- Up to 52 Mbps *downlink* (router-to-home)



Home network: Broadband

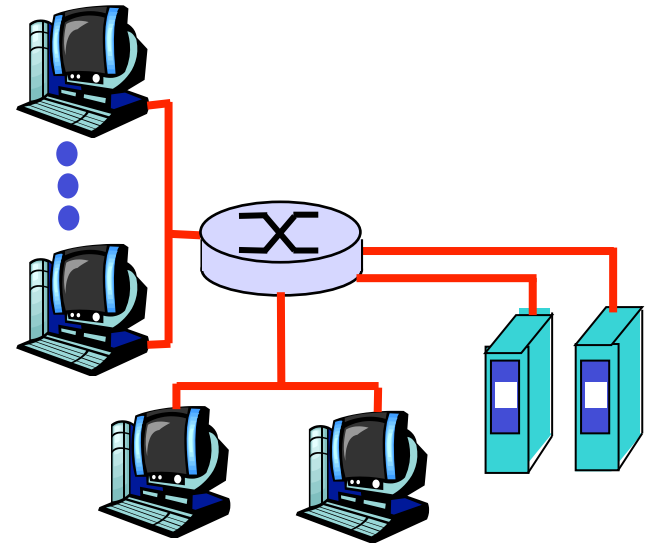
An example

- ❑ **HFC:** hybrid fiber coax
 - Asymmetrical: e.g. 25 Mbps downlink, 5 Mbps uplink
- ❑ **Network** of copper cable and optical fiber connects homes to ISP routers
 - Shared access to router for several homes
 - Problems: congestion, dimensioning



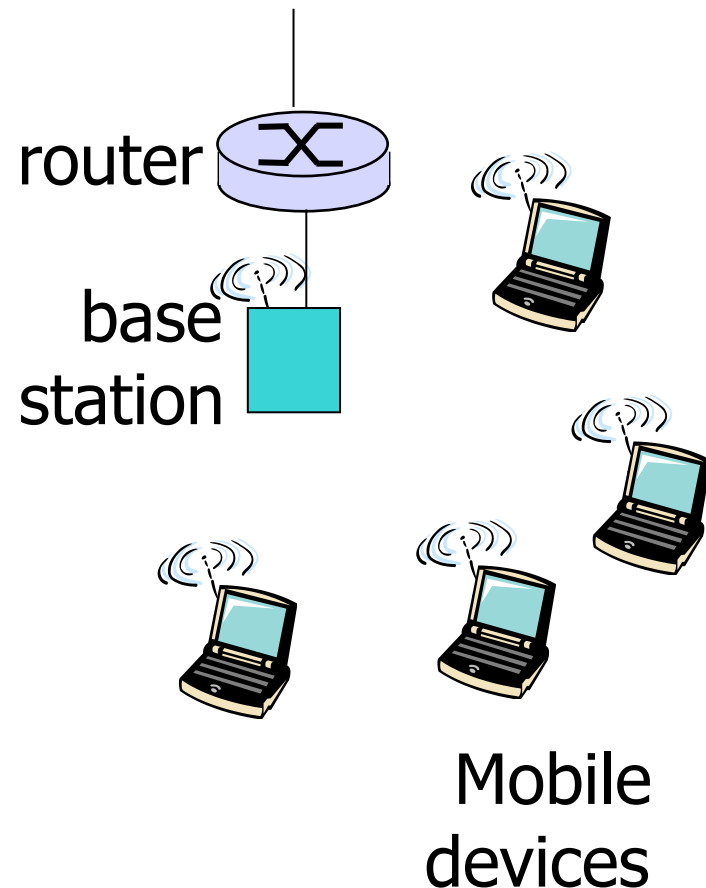
Institutional access networks (LAN)

- ❑ Company/university **local area network (LAN)** connects end systems to the rest of the net
- ❑ **Ethernet:**
 - Shared or dedicated cable connects end systems and routers
 - 10 Mbps, 100Mbps, Gigabit Ethernet



Wireless access networks

- ❑ Shared wireless access networks connect end systems to routers
- ❑ **Wireless LANs:**
 - radio spectrum replaces cable
 - E.g.
 - IEEE 802.11g - 54 Mbps
 - IEEE 802.11h – 100 Mbps
 - IEEE 802.11n – 600Mbps
 - IEEE 802.11ac – 1,3Gbps
- ❑ **Wireless access over long distances**
 - 3G/4G for example...



Physical medium

- ❑ **Physical link:** a sent bit propagates through the link
- ❑ **Closed media:**
 - Signals propagate in cable media (copper, fiber)
- ❑ **Open media:**
 - Signals propagate freely, e.g. radio.

Twisted Pair (TP)

- ❑ Two isolated copper cables
 - Category 3: traditional telephone cables, 10 Mbps Ethernet
 - Category 5 TP: 100Mbps Ethernet
 - Category 6 TP: 1Gpbs Ethernet



Physical medium: coax, fiber

Coaxial cable

- ❑ Wire (signal carrier) in a wire (shielding)
 - baseband: a single channel on a cable
 - broadband: multiple channels on a cable
- ❑ Bi-directional
- ❑ Typically used for 100Mbps Ethernet.



Fiber optic cable

- ❑ Optical fiber that carries light impulses
- ❑ High-speed transfer:
 - High-speed point-to-point transmission
- ❑ Low error rate
- ❑ Longer distances
 - 100Mbps, 1-100Gbps Ethernet

Physical media: radio

Radio

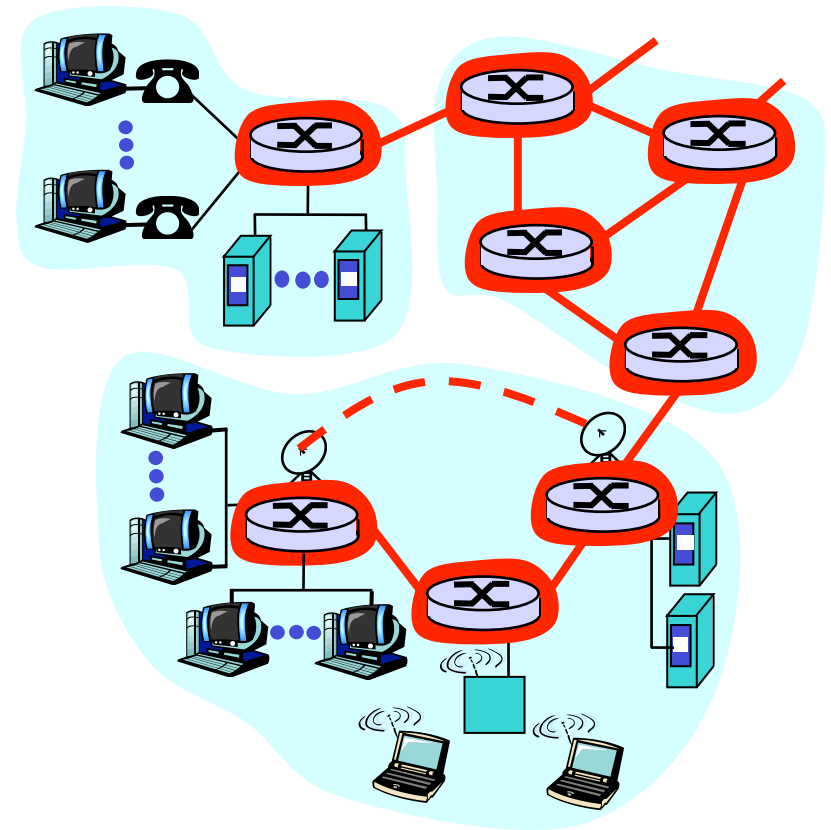
- ❑ Signal in electromagnetic spectrum
- ❑ No physical "cable"
- ❑ Bi-directional
- ❑ Effects of environment on the distribution:
 - Reflection
 - Obstruction by blocking objects
 - Interferences

Types of radio links

- ❑ **microwaves**
 - E.g. up to 45 Mbps
- ❑ **WLAN**
 - 54Mbps, 600Mbps, 1,3Gbps
- ❑ **wide-area**
 - 3G, 14,4Mbps (in theory)
- ❑ **satellite**
 - Up to 50Mbps per channel (or several thinner channels)
 - 270 ms end-to-end delay (limited by speed of light).

Core networks

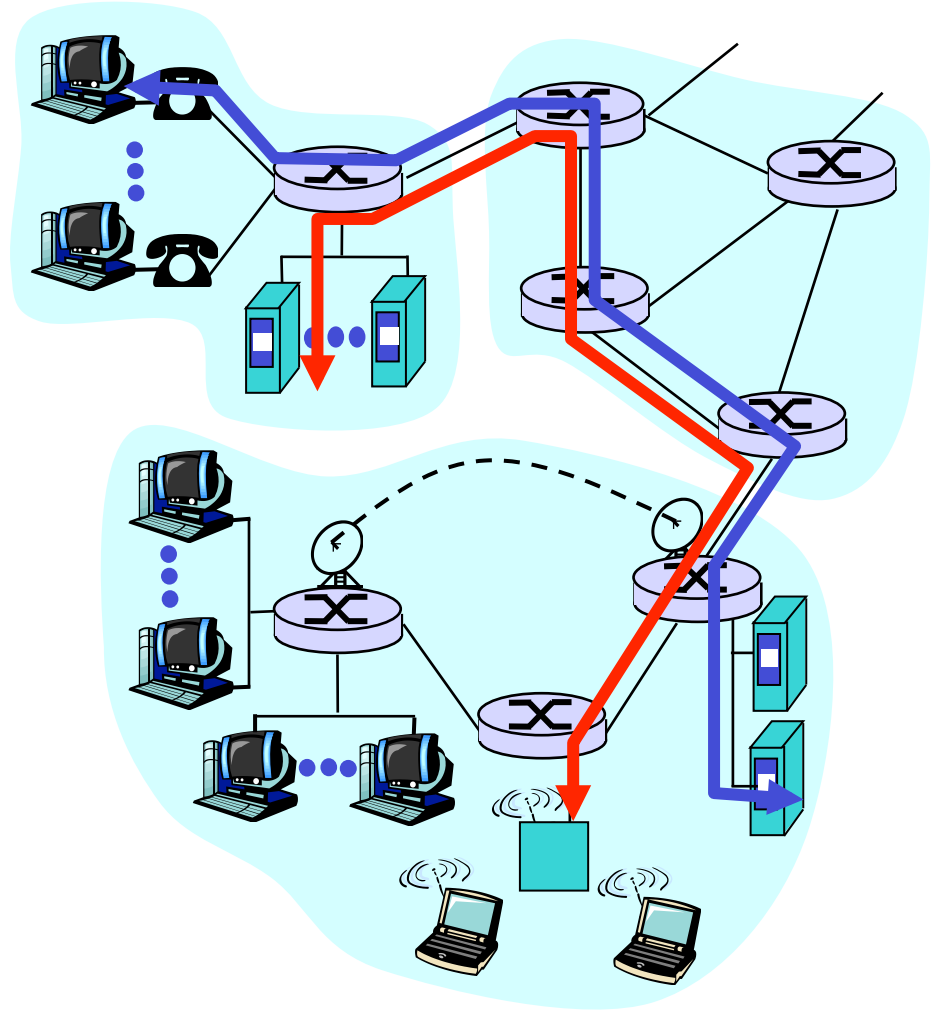
- ❑ Graph of interconnected routers
- ❑ One fundamental question: how is data passed through the net?
 - Circuit switching
 - Packet switching
- ❑ **Circuit switching**
 - Dedicated line through the network
- ❑ **Packet switching**
 - Discrete *data units* are sent through the network



Core networks: Circuit Switching

End-to-end resource reservation for a "session"

- ❑ Setup phase is required
- ❑ Dedicated resources (no sharing)
- ❑ Link bandwidth, router capacity
- ❑ Guaranteed throughput



Core networks: Circuit Switching

Historical:

- ❑ Analog telephone networks
- ❑ Network consists of resources
 - Cables
 - Switches with relays
- ❑ Establish a physical connection
 - Relays switch to connect cables physically
 - Create a circuit
 - Guaranteed resources
 - No difference between talking and silence

Modern:

- ❑ Networks consist of resources
 - Cables
 - Routers or switches
 - Network resources **can be shared**
- ❑ Establish a connection
 - Switches reserve part of available resource
- ❑ Division of link bandwidth into parts
 - Frequency division
 - Time division

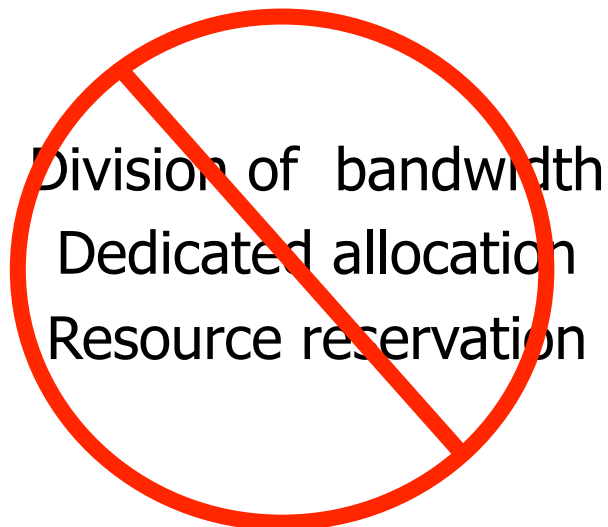
Core networks: Packet Switching

Each end-to-end data stream is
divided into packets

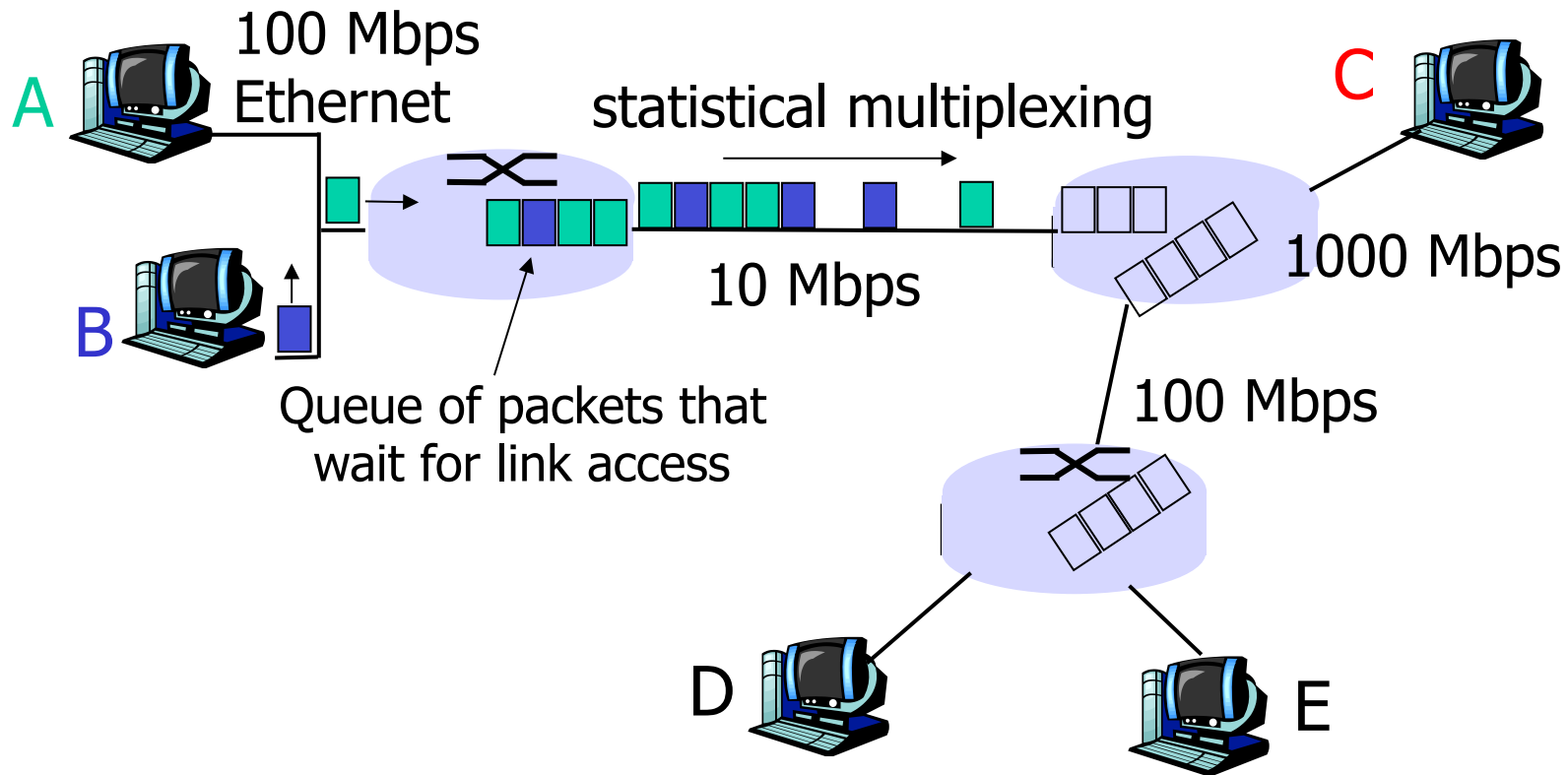
- ❑ Data streams *share* network resources
- ❑ Each packet uses the entire bandwidth of a link
- ❑ Resources are used as needed

Competition for resources:

- ❑ Combined resource need can exceed the available resources
- ❑ Congestion: packets are queued in front of “thin” links
- ❑ Store and forward: packets move one link at a time
 - Send over a link
 - Wait for your turn at the next link



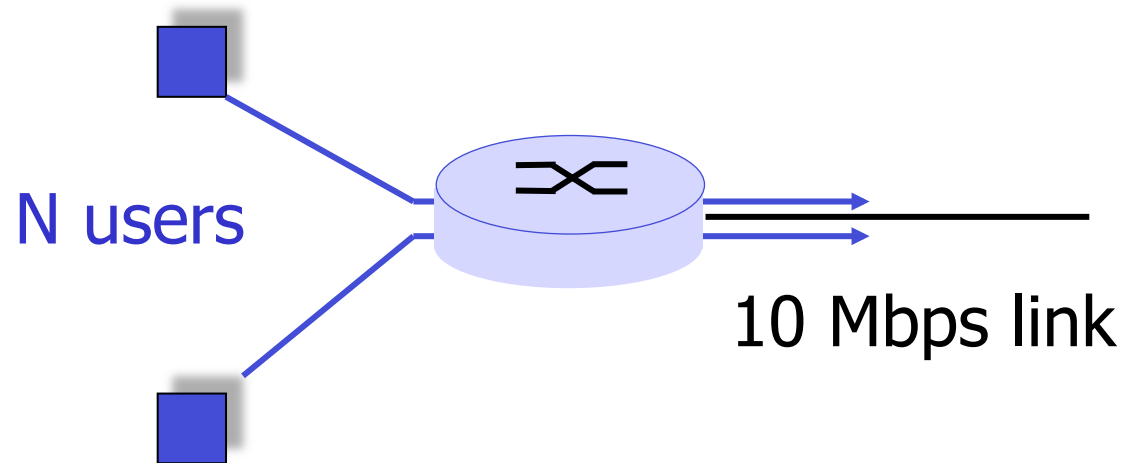
Core networks: Packet switching



Packet switching versus circuit switching

Packet switching allows more users in the net!

- ❑ 10 Mbps link
- ❑ Each user
 - 1Mbps when “active”
 - Active 10% of the time, at random times
- ❑ Circuit switching
 - max 10 users
 - Loss probability: 0%
 - Waste: ~90% capacity
- ❑ Packet switching
 - >10 may be active concurrently!
 - Loss probability >0%
 - Waste: < 90% capacity



Packet switching versus circuit switching

Is packet switching always the best approach?

- ❑ Good for data with “bursty” behavior
 - Resource sharing
 - No “setup phase” required
- ❑ In a congested network: delay and packet loss
 - Protocols/mechanisms required for reliable traffic and congestion control
- ❑ How to achieve a behavior like that of circuit switching?
 - Bandwidth guarantees are required for audio/video applications

QoS concepts have to be used for that purpose!

Delay in packet switching networks

Packet experiences **delay** on the way from sender to receiver

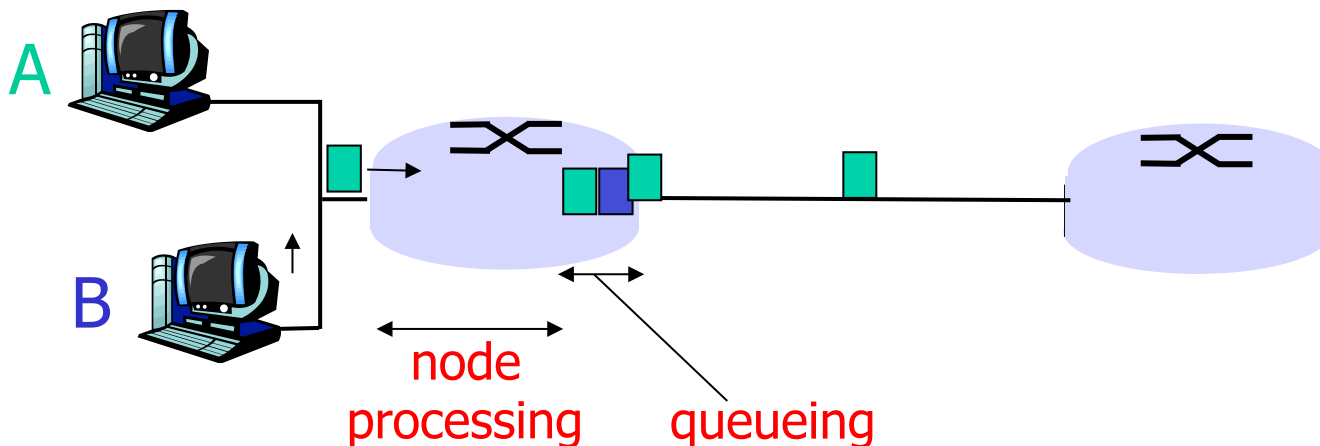
□ **four** sources of delay in each hop.

□ Node processing:

- Determining the output link – address lookup
- Checking for bit errors

□ Queuing

- Waiting for access to the output link
- Depends on the congestion level of the router



Delay in packet switching networks

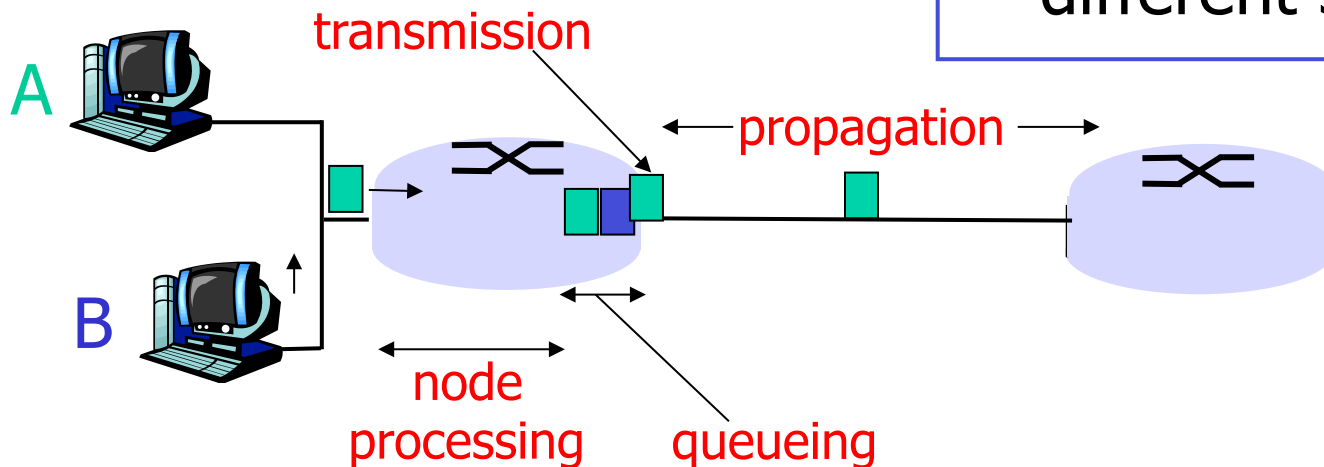
Transmission delay:

- ❑ R = link bandwidth (bps)
- ❑ L = packet size (bits)
- ❑ Time required to send a packet onto the link = L/R

Propagation delay:

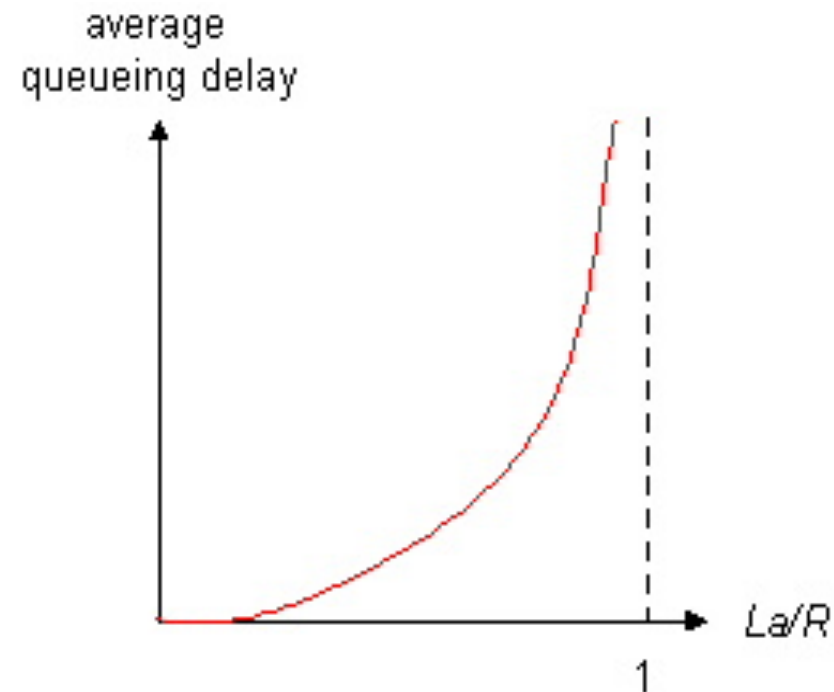
- ❑ d = physical link length (m)
- ❑ s = propagation speed in the medium ($\sim 2 \times 10^8$ m/sec)
- ❑ Propagation delay = d/s

Note: s and R are of very different size!



More about queueing delays

- ❑ R = link bandwidth (bps)
 - ❑ L = packet length (bits)
 - ❑ a = average packet arrival rate
- traffic intensity = La/R

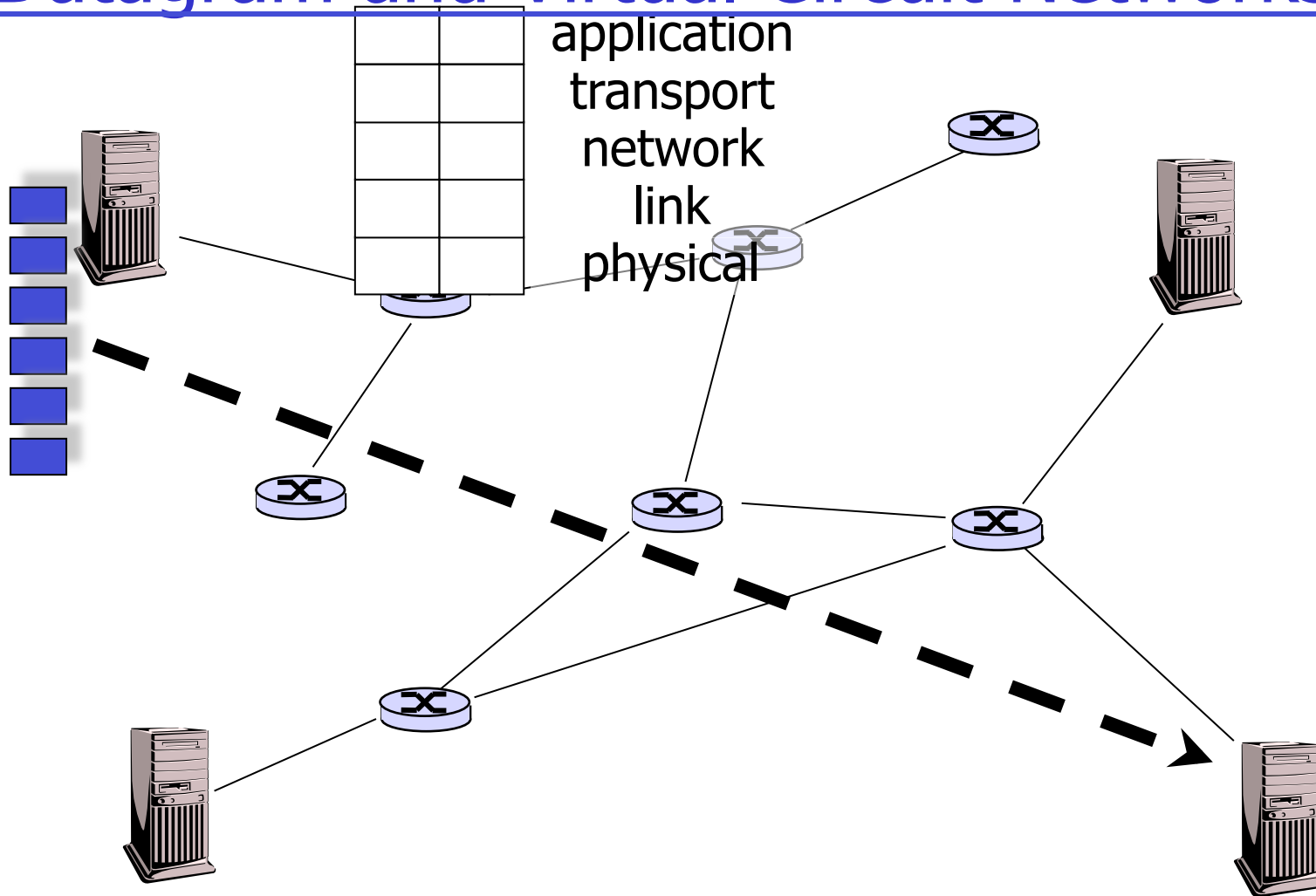


- ❑ $La/R \sim 0$: average queuing delay is small
- ❑ $La/R \rightarrow 1$: queuing delay grows
- ❑ $La/R > 1$: more data is arriving at the link than it can handle \rightarrow link goes into congestion (Average delay is infinite!)

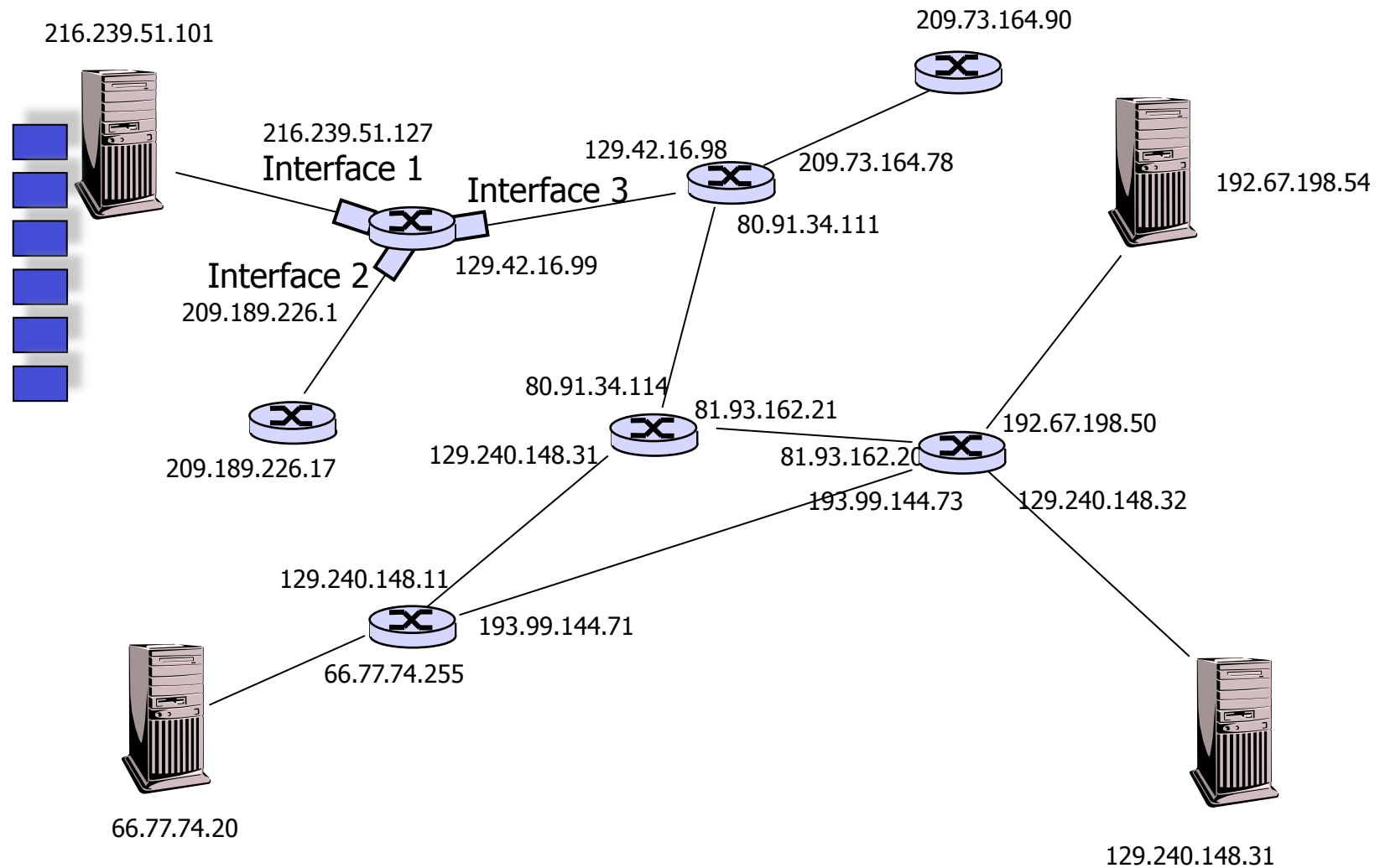
Packet switched network: Routing

- ❑ Goal: move packets from router to router between source and destination
 - There are two methods to find the path of packets.
- ❑ **Datagram network:**
 - *Destination address* determines the next hop.
 - Path can change during the sessions.
 - Routers need no information about sessions.
 - Analogy: ask for the way while you drive.
- ❑ **Virtual circuit network:**
 - Each packet has a “tag” (virtual circuit ID), which determines the next hop.
 - Path is determined when connection is set up, and remains the same for the entire session.
 - Routers need state information for each virtual circuit.

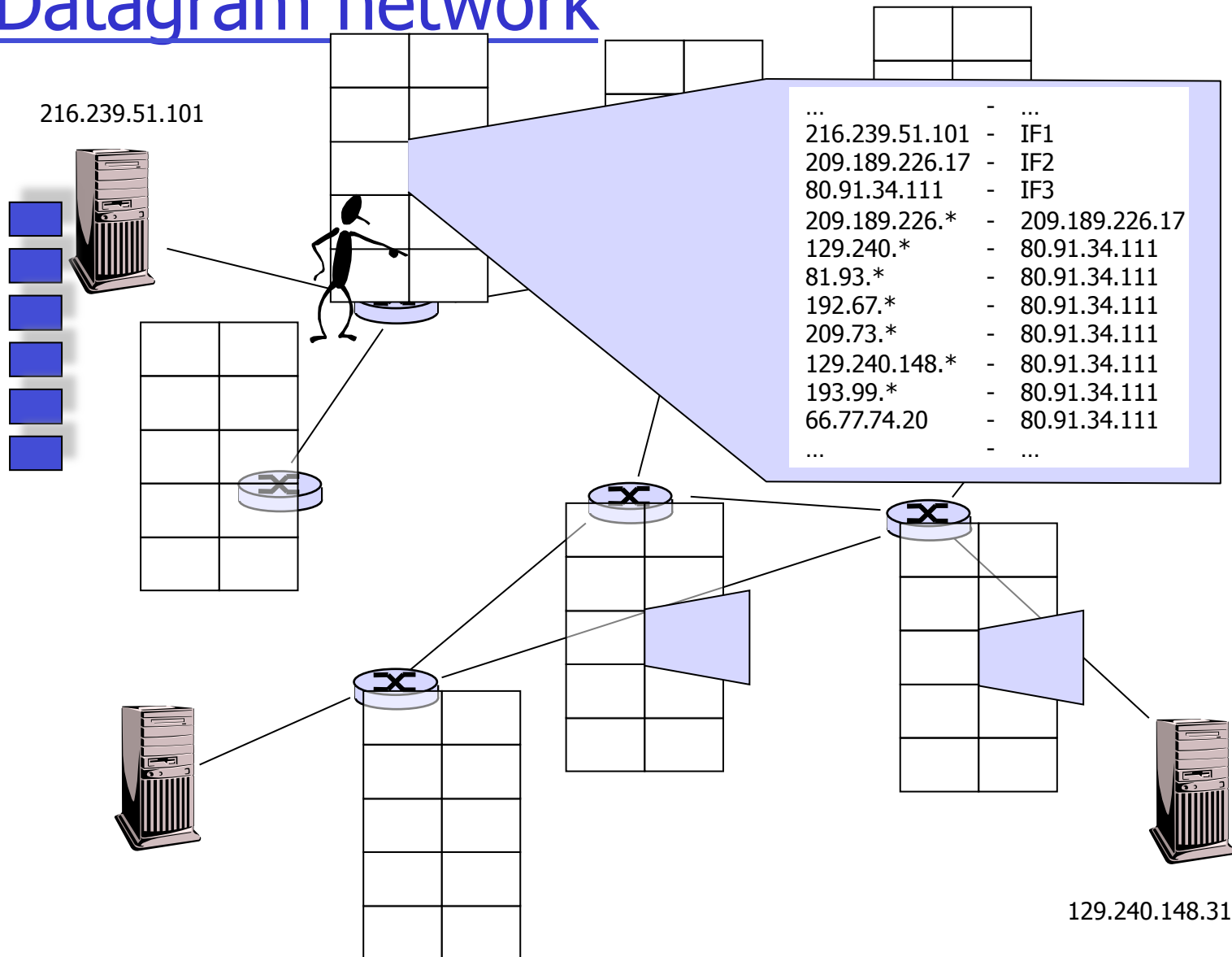
Datagram and Virtual Circuit Networks



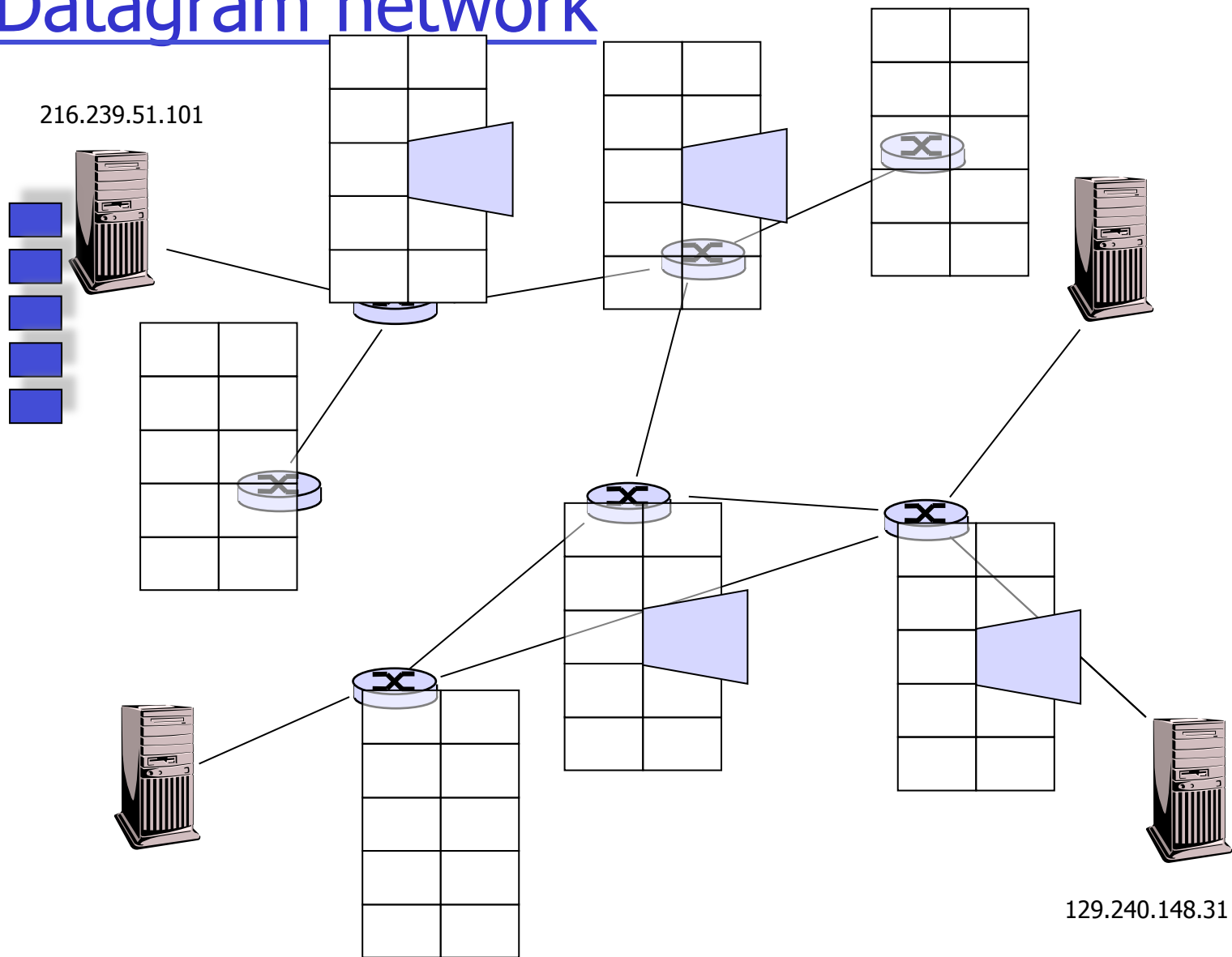
Datagram and Virtual Circuit Networks



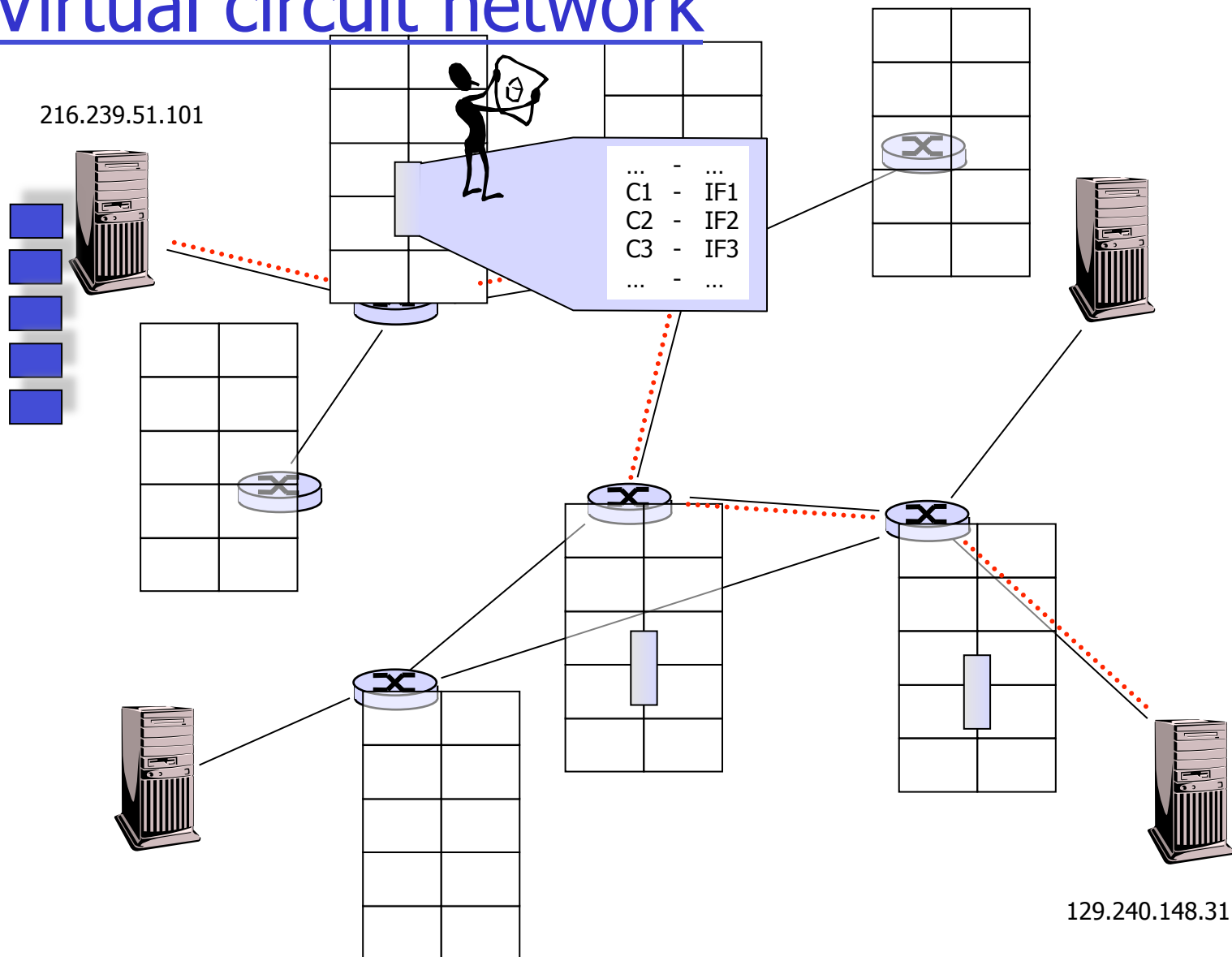
Datagram network



Datagram network



Virtual circuit network



Network layer: IP

Datagram switching

- ❑ IP
 - Internet Protocol
 - Datagram service of the Internet
 - RFC 791

IP offers:

- ❑ Addressing
- ❑ Routing
- ❑ Datagram service
 - Unreliable
 - Unordered

IP networks can use virtual circuits

- ❑ IPv4: circuit is one hop
- ❑ IPv6: can have a tag

Connection-oriented service

Goal: data transfer between end systems

□ Start of communication

- *Handshaking*
- Initial preparation of data transfer
- Hi!, hi! Is a human handshaking protocol
- Creates a "*state*" in the two machines that communicate.
- End systems know their communication partners

□ During communication

- *Connection*
- End system expects messages from connected end system
- End system knows when messages belong to the connection

□ End of communication

- *Teardown*
- Bye! Bye! Is a human teardown protocol
- New handshake required for re-establishing connection

Connectionless service

Goal: data transfer between end systems

❑ As before!

❑ Start of communication

- *No connection setup*
- No preparation for data transfer
- Programs must expect messages at all times

❑ During communication

- *No connection*
- No state in the machines
- Senders don't know whether messages are expected
- Sender must identify itself in each message

❑ End of communication

- *No teardown*
- Just stop sending

Services over Switching Approaches

- ❑ Services requested
 - Between end systems
 - Connection-oriented service
 - Connectionless service
- ❑ Switching approaches
 - From host to host
 - Circuit switching
 - Packet switching

	Connection-oriented service	Connection-less service
Circuit switching	Fits well	Setup wasted
Packet switching	Additional work needed	Fits well

Transport layer: TCP

Connection-oriented service

- ❑ TCP
 - Transmission Control Protocol
 - Connection-oriented service of the Internet
 - RFC 793

TCP offers:

- ❑ *Connections*
 - *Handshake, end-system state, teardown*
- ❑ *Reliable, ordered, stream-oriented data transfer*
 - Loss: acknowledgements and retransmissions
- ❑ *Flow control:*
 - Send not faster than receiver can receive
- ❑ *Congestion control:*
 - Send slower when the network is congested.

Transport layer: UDP

Connectionless service

- ❑ UDP
 - User Datagram Protocol
 - Connectionless service of the Internet
 - RFC 768

UDP offers:

- ❑ *No connections*
 - Send immediately
- ❑ *Unreliable, unordered, packet-oriented data transfer*
 - Loss: messages are simply lost
 - Messages arrive exactly as send
- ❑ *No flow control*
 - Send as fast as programs want to
- ❑ *No congestion control*
 - Ignore network problems

Transport layer: applications

Applications that use TCP:

- ❑ HTTP (WWW)
- ❑ FTP (file transfer)
- ❑ SMTP (email)

- ❑ Telnet (remote login)

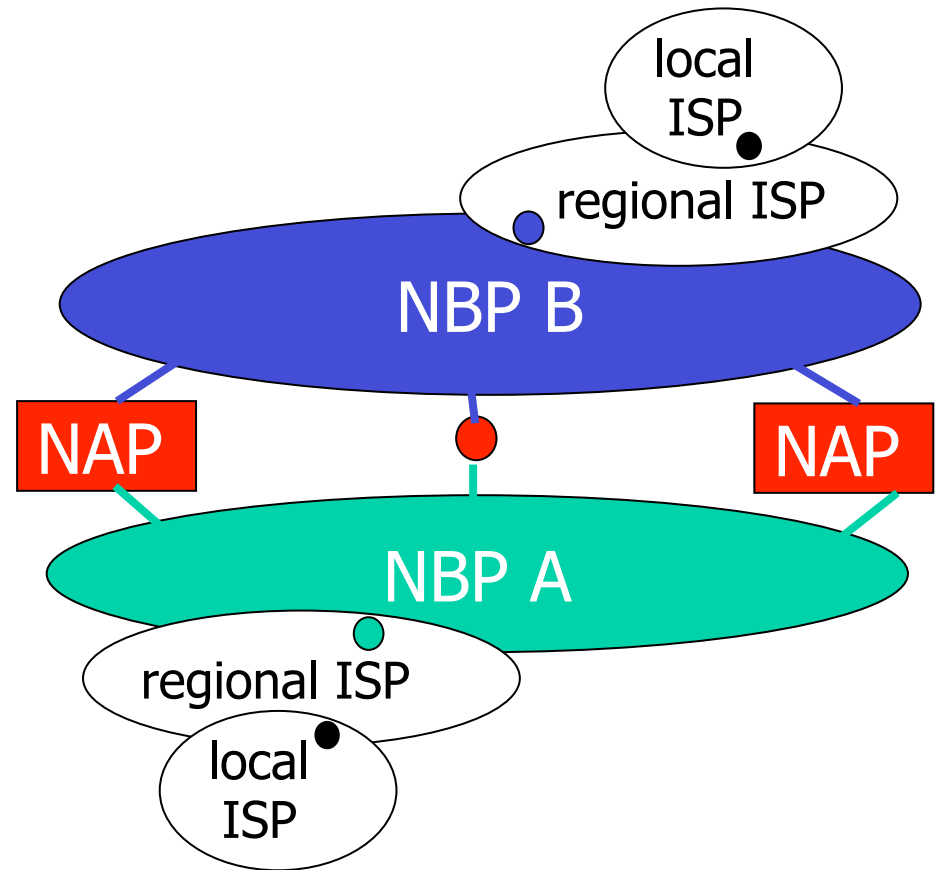
Applications that use UDP:

- ❑ Streaming media
- ❑ Video conferencing
- ❑ Internet telephony

- ❑ NTP (network time protocol)

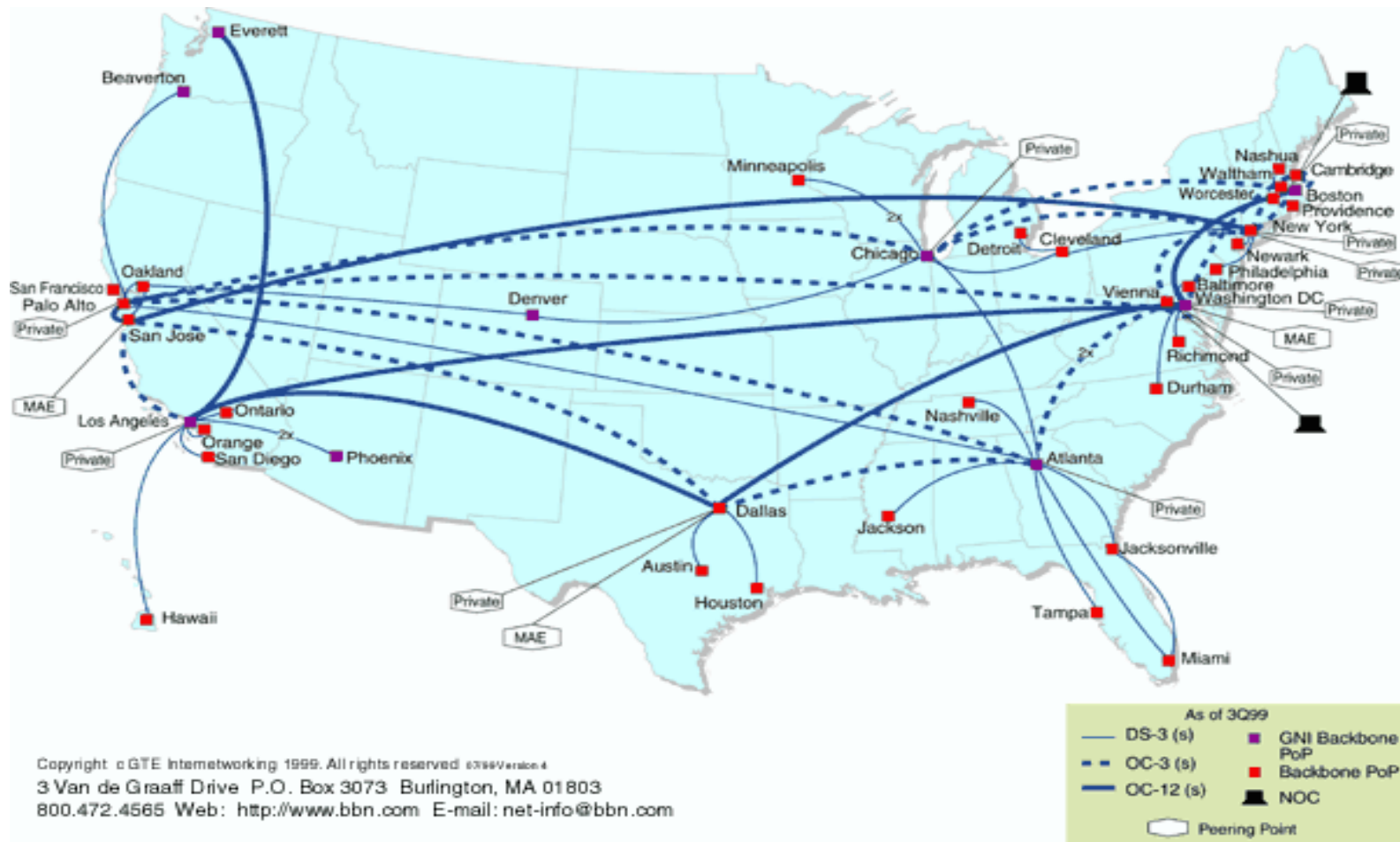
Internet structure: network of networks

- ❑ More or less hierarchical
- ❑ **National/international “backbone providers” (NBPs)**
 - These interconnect either privately, or at so-called Network Access Point (NAPs)
- ❑ **regional ISPs**
 - Connect to NBPs
- ❑ **local ISPs, companies**
 - Connect to regional ISPs



National Backbone Provider

example BBN/GTE US backbone network



History of the Internet

1961-1972: Early packet-switching concepts

- ❑ 1961: Kleinrock – queueing theory proves that packet switching is effective
- ❑ 1964: Baran – packet switching in military networks
- ❑ 1967: ARPAnet starts Advanced Research Projects Agency
- ❑ 1969: first ARPAnet node operational
- ❑ 1972:
 - ARPAnet publically demonstrated
 - NCP (Network Control Protocol) first machine-machine protocol
 - first e-mail program
 - ARPAnet has 15 nodes

History of the Internet

1972-1980: Internetworking – new, proprietary networks

- ❑ 1970: ALOHAnet satellite network on Hawaii
- ❑ 1973: Metcalfe's doctor thesis proposes Ethernet
- ❑ 1974: Cerf and Kahn – architecture to the interconnection of many networks
- ❑ End of the 70s: proprietary architectures: DECnet, SNA, XNA
- ❑ 1979: ARPAnet has 200 nodes

Cerf og Kahn's internetworking principles:

- Minimalism, autonomy – no internal network changes necessary to interconnect networks
- best effort service model
- Stateless routers
- Decentralized control

This defines mostly today's Internet architecture

History of the Internet

1980-1990: new protocols – the Net grows

- ❑ 1983: first use of TCP/IP
- ❑ 1982: e-mail protocol SMTP defined
- ❑ 1983: DNS defined to translate a name into an IP address
- ❑ 1985: ftp protocol defined
- ❑ 1988: TCP congestion control
- ❑ New national networks: Csnet, BITnet, NSFnet, Minitel
- ❑ 100,000 machine connected to the Net.

History of the Internet

1990's: commercialization, WWW

□ Early 1990s: WWW

- hypertext [Bush 1945, Nelson 1960's]
- HTML, http: Berners-Lee
- 1994: Mosaic, later Netscape
- late 1990s: commercialization of the WWW

Late 1990s:

- ca. 50 million machines on the Internet
- ca. 100 million+ users
- backbone links operate at 1 Gbps

Summary

Covering a large area!

- ❑ Overview over the Internet
- ❑ What is a protocol?
- ❑ Network components
- ❑ Throughput, loss, delay
- ❑ Layering and service models
- ❑ backbone, NAPer, ISPer

Hopefully you have now:

- ❑ An impression and overview of the area
- ❑ More depth and details in the following lessons, and in later courses