# inf2080 oppgave 4.22 

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Given two disjoint co-Turing-recognizable languages there exists a decidable language separating them.

Let $A$ and $B$ be two disjoint co-Turing-recognizable languages so that $A \cap B=\emptyset$, $\operatorname{coT} M_{A}$ recognizes $\bar{A}$ and $\operatorname{coT} M_{B}$ recognizes $\bar{B}$. We construct $T M_{C}$ that decides $C$, a language separating $A$ and $B$.
$T M_{C}=$ "on input $\omega$ :

1. Simulate running $\operatorname{coT} M_{A}$ on $\omega$ and $\operatorname{coT} M_{B}$ on $\omega$ in parallel (alternating between $\operatorname{coTM} M_{A}$ and $\operatorname{coTM} M_{B}$ ).
2. If at any time $\operatorname{coT} M_{A}$ accepts, $R E J E C T$. If at any time $c o T M_{B}$ accepts, ACCEPT."
$C$ satisfies the criteria for separating $A$ and $B$ :

- $A \subseteq C$ : On input $\omega \in A \operatorname{coTM}_{A}$ will loop or $R E J E C T$, $\operatorname{coTM}_{B}$ will $A C C E P T$ since $\omega \notin B$ and $T M_{C} A C C E P T$ s.
- $B \subseteq \bar{C}$ : On input $\omega \in B \operatorname{coT} M_{B}$ will loop or REJECT, coTM $A$ will $A C C E P T$ since $\omega \notin A$ and $T M_{C} R E J E C T \mathrm{~s}$.
$T M_{C}$ is a decider: if $\omega \notin A \cup B$ then $\operatorname{coTM}_{A}$ and $\operatorname{coT} M_{B}$ will race. Both machines eventually accept, but the machine that finishes computation first decides whether $T M_{C}$ accepts or rejects.

