

# INF3100: Databasesystemer

## Oppgavesett 12

**Oppgave 20.5.2:** In this exercise, we need a notation for describing sequences of messages that can take place during a two-phase commit. Let  $(i, j, M)$  mean that site  $i$  sends the message  $M$  to site  $j$ , where the value of  $M$  and its meaning can be  $P$  (prepare),  $R$  (ready),  $D$  (don't commit),  $C$  (commit), or  $A$  (abort). We shall discuss a simple situation in which site 0 is the coordinator, but not otherwise part of the transaction, and sites 1 and 2 are the components. For instance, the following is one possible sequence of messages that could take place during a successful commit of the transaction:

$(0, 1, P), (0, 2, P), (2, 0, R), (1, 0, R), (0, 2, C), (0, 1, C)$

a) Give an example of a sequence of messages that could occur if site 1 wants to commit and site 2 wants to abort.

**Oppgave 20.7.1:** Given the circle of nodes of Fig. 20.14, where do key-value pairs reside if the key hashes to: (a) 35 (b) 20 (c) 60?

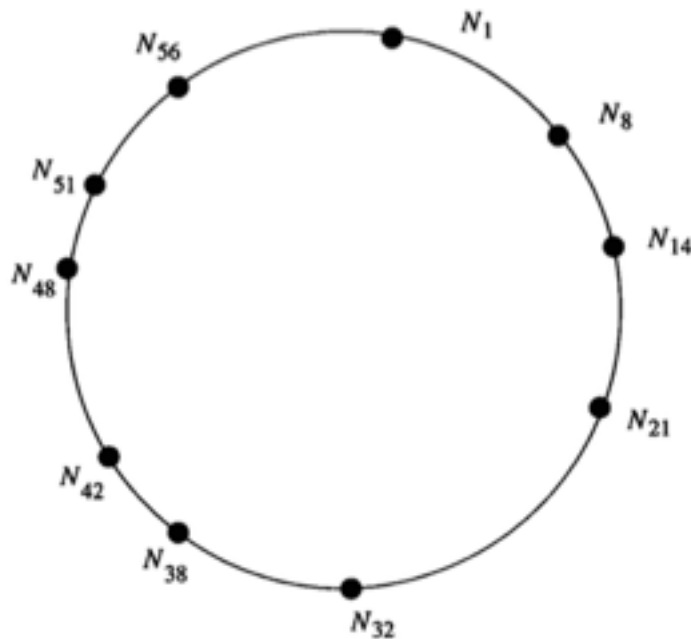


Figure 20.14: A chord circle

To place a node in the circle, we hash its ID  $i$ , and place it at position  $h(i)$ . We shall henceforth refer to this node as  $N(hi)$ . Thus, for example, in Fig. 20.14,  $N21$  is a node whose ID  $i$  has  $h(i) = 21$ . The successor of each node is the next higher one clockwise around the circle. For example, the successor of  $N21$  is  $N32$ , and  $N1$  is the successor of  $N56$ . Likewise,  $N21$  is the predecessor of  $N32$ , and  $N56$  is the predecessor of  $N1$ .

The nodes are located around the circle using a hash function  $h$  that is capable of mapping both keys and node ID's (e.g., IP-addresses) to  $m$ -bit numbers, for some  $m$ . In Fig. 20.14, we suppose that  $m = 6$ , so there are 64 different possible locations for nodes around the circle. In a real application,  $m$  would be much larger.

Key-value pairs are also distributed around the circle using the hash function  $h$ . If  $(K, V)$  is a key-value pair, then we compute  $h(K)$  and place  $(K, V)$  at the lowest numbered node  $N_j$  such that  $h(K) < j$ . As a special case, if  $h(K)$  is above the highest-numbered node, then it is assigned to the lowest-numbered node. That is, key  $K$  goes to the first node at or clockwise of the position  $h(K)$  in the circle.

**Oppgave 20.7.2:** Given the circle of nodes of Fig. 20.14, construct the finger tables for: (a)  $N14$  (b)  $N51$

**Oppgave 20.7.3:** Given the circle of nodes of Fig. 20.14, what is the sequence of messages sent if:

- a)  $N14$  searches for a key that hashes to 27
- b)  $N8$  searches for a key that hashes to 5
- c)  $N56$  searches for a key that hashes to 54

**Oppgave 16.6.4:** Consider the join of relations  $R(a,b)$ ,  $S(b,c)$ ,  $T(c,d)$ , and  $U(a,d)$ , where  $R$  and  $U$  each have 1000 tuples, while  $S$  and  $T$  each have 200 tuples. Further, there are 200 values of all attributes of all relations, except for attribute  $c$ , where  $V(S,c) = V(T,c) = 20$ .

- a) What is the order selected by the greedy algorithm? What is its cost?
- b) What is the optimum join ordering and its cost?