

## INF3100: Databasesystemer – Oppgavesett 11

**Oppgave 16.3.2:** Replace the natural joins in the following expression by equivalent theta-joins and projections. Tell whether the resulting theta-joins form a commutative and associative group.

- a)  $(R(a, b) \bowtie S(b, c)) \bowtie (T(c, d) \bowtie U(d, e))$
- b)  $(R(a, b) \bowtie S(b, c)) \bowtie (T(c, d) \bowtie U(a, d))$
- c)  $(R(a, b) \bowtie S(b, c)) \bowtie_{R.a < T.c} T(c, d)$

**Oppgave 16.6.4:** Consider the join of relations  $R(a, b)$ ,  $S(b, c)$ ,  $T(c, d)$ , and  $U(a, d)$ , where  $R$  and  $U$  each have 1000 tuples, while  $S$  and  $T$  each have 200 tuples. Further, there are 200 values of all attributes of all relations, except for attribute  $c$ , where  $V(S, c) = V(T, c) = 20$ .

- a) What is the order selected by the greedy algorithm? What is its cost?
- b) What is the optimum join ordering and its cost?

**Oppgave 16.x.2:** Gitt en hotelldatabase med relasjonene

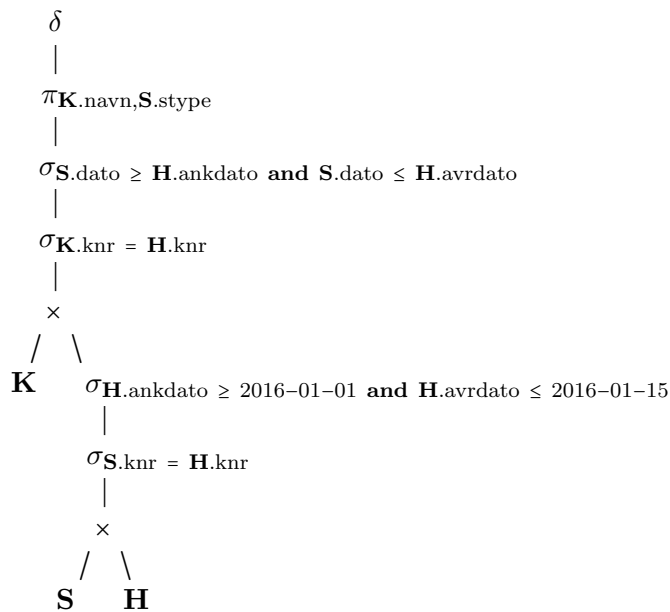
**Kunde**(knr, navn, adr)

**Hotellrom**(knr, ankdato, avrdato, rom, pris)

**Service**(knr, dato, stype, pris)

**Kunde** og **Hotellrom** burde være rimelig selvforklarende. **Kunde** har primærnøkkel (knr) og **Hotellrom** primærnøkkel (knr, ankdato). **Service** inneholder informasjon om servicetilbud som hotellgjestene har benyttet seg av, som spa, frisør, restaurantbesøk mm., med prisinformasjon (som blir lagt til den endelige hotellregningen ved avreise). **Service** har primærnøkkel (knr, dato, stype).

I relasjonsalgebrauttrykket under er **K**, **H** og **S** forkortelser for henholdsvis **Kunde**, **Hotellrom** og **Service**.



- a) Hva uttrykker spørringen?
- b) Optimaliser uttrykket.

**Oppgave 15.2.3:** Give one-pass algorithms for each of the following join-like operators:

- a)  $R \overset{\circ}{\bowtie}_L S$ , assuming  $R$  fits in memory (see Section 5.2.7 for definitions involving outerjoins).
- b)  $R \overset{\circ}{\bowtie}_L S$ , assuming  $S$  fits in memory.

**Oppgave 15.6.1:** Suppose  $B(R) = 10,000$  and  $T(R) = 500,000$ . Let there be an index on  $R.a$ , and let  $V(R, a) = k$  for some number  $k$ . Give the cost of  $\sigma_{a=0}(R)$  as a function of  $k$ , under the following circumstances. You may neglect disk I/O's needed to access the index itself.

- a) The index is not clustering.
- b) The index is clustering.
- c)  $R$  is clustered, and the index is not used.

**Oppgave 20.5.2:** In this exercise, we need a notation for describing sequences of messages that can take place during a two-phase commit. Let  $(i, j, M)$  mean that site  $i$  sends the message  $M$  to site  $j$ , where the value of  $M$  and its meaning can be  $P$  (prepare),  $R$  (ready),  $D$  (don't

commit),  $C$  (commit), or  $A$  (abort). We shall discuss a simple situation in which site 0 is the coordinator, but not otherwise part of the transaction, and sites 1 and 2 are the components. For instance, the following is one possible sequence of messages that could take place during a successful commit of the transaction:

$$(0, 1, P), (0, 2, P), (2, 0, R), (1, 0, R), (0, 2, C), (0, 1, C)$$

- a) Give an example of a sequence of messages that could occur if site 1 wants to commit and site 2 wants to abort.