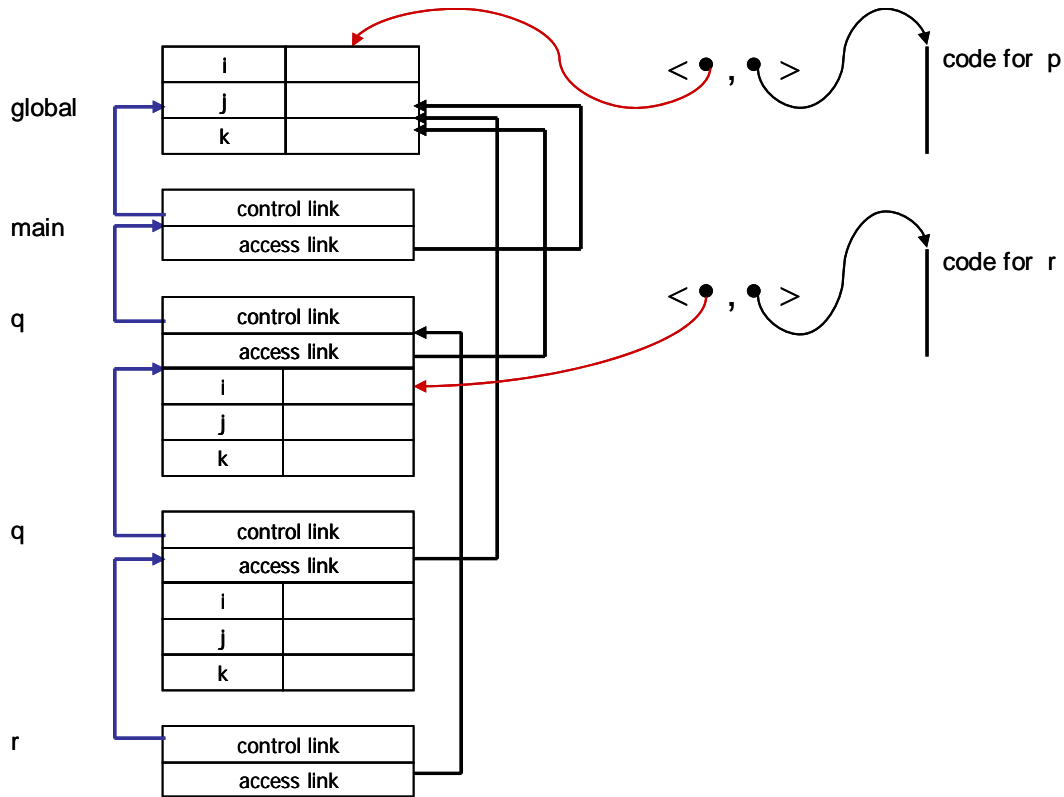


Oppgave 1 Runtime-systemer, skoping, typer (vekt 40%)

1a

| | statisk scoping | dynamisk scoping |
|-----------------------------|-----------------|------------------|
| 1. utførelse av en print(k) | 3 | 9 |
| 2. utførelse av en print(k) | 6 | 9 |
| 3. utførelse av en print(k) | 9 | 9 |
| 4. utførelse av en print(k) | 6 | 9 |

1 b & c



1d

1)

If $dist' < dist$, then it should be possible to substitute a call $dist'(somePoint)$ with a call $dist'(somePoint)$. $dist'$ will be able to manipulate properties of $ColorPoint$, that an object of class $Point$ will not have. We therefore have to check at run-time if $somePoint$ points to a $Point$ or $ColorPoint$ object.

However, with $dist < dist'$ a call $dist'(someColorPoint)$ may be substituted with the call $dist(someColorPoint)$, and as $dist$ can only use properties defined for $Point$, and actual parameters must be of class $ColorPoint$ they will always have $Point$ properties.

1e

2)

If `somePoint < somePoint'`, then it should be possible to substitute a call `somePoint'()` with a call `somePoint()`. Such a call may be:

```
aColorPoint = somePoint'(1)
```

If `somePoint'(1)` is substituted with `somePoint(1)`, then one will try to assign a `Point` object to the reference `aColorPoint` which is typed with `ColorPoint`, and this will require a run-time type check.

However, if `somePoint' < somePoint`, then a call `somePoint()` may just as well be substituted with a call `somePoint'()`. Such a call may be:

```
aPoint = somePoint(1)
```

If `somePoint(1)` is substituted with `somePoint'(1)`, then one will try to assign a `ColorPoint` object to the reference `aPoint` typed by `Point`, and this will always be legal.

Oppgave 2 ML (vekt 40%)**2a**

```

fun f(g,h) = g(g(h)) * 3;

0. w = int -> int -> int
1. t = r -> s
2. t = s -> b
3. w = b -> c
4. c = int -> d
5. e = t * r -> d

From 1 & 2:
6. r = s = b

From 3 & 4:
7. w = b -> int -> d

From 0 & 7:
8. b = int
9. d = int

From 6 & 8:
10. r = int
11. s = int

From 1, 10 & 11:
12. t = int -> int

From 9, 10 & 12:
13. e = (int -> int) * int -> int

```

2b

```

1) fun firstElems xs = map hd xs;
2) fun prodFirstElems xs = foldl op* 1 (firstElems xs);
   Another solution:
   fun prodFirstElems2 xs = foldr op* 1 (firstElems xs);

```

2c

ONE SOLUTION:

```

fun f(0,count) = count
  | f(1,count) = 0.0/0.0
  | f(3,count) = ~1.0
  | f(x,count) = f(x-2, count+1.0);

```

ANOTHER POSSIBLE SOLUTION:

```

fun fAux(0,count, flag) = count
  | fAux(1,count, flag) = if flag=1 then 0.0/0.0 else ~1.0
  | fAux(x,count, flag) = fAux(x-2, count+1.0, flag)

fun eqTof(1,count) = fAux(1, count, 1)
  | eqTof(x,count) = fAux(x, count, 0);

```

Oppgave 3 Prolog (vekt 20%)

3a

a)

For example:

```
mother(sarah,per).
mother(sarah,anne).
mother(anne,sofia).
mother(sofia,carlos).
```

b)

```
son(X,Y) :- father(Y,X), male(X).

daughter(X,Y) :- father(Y,X), female(X).

brother(X,Y) :- father(Z,X), father(Z,Y), X\==Y.

parent(X,Y) :- father(x,Y).
parent(X,Y) :- mother(x,Y).

uncle(Uncle,Person) :- brother(Uncle,Parent), parent(Parent,Person).
```

c)

- All the persons who have at least one brother.

Define first:

```
atleastonebrother(X) :- (son(X,Y); daughter(X,Y)), son(Z,Y), X\==Z.
    % Alternative: atleastonebrother(X) :- brother(X,Y), male(Y).

atleastonebrother(X).
```

- All the persons who are uncle of a female

Define first:

```
unclefemale(X) :- uncle(X,Y), female(Y).

unclefemale(X).
```

3b

Natural numbers may be defined as follows in Prolog:

```
natural_number(0).
natural_number(s(X)) :- natural_number(X).

plus(0,X,X) :- natural_number(X).
plus(s(X),Y,s(Z)) :- plus(X,Y,Z).
```

```
prod(0,X,0) :- natural_number(0).  
prod(s(X),Y,Z) :- prod(X,Y,XY), plus(XY,Y,Z).  
  
exp(s(N),0,0).  
exp(0,s(X),s(0)).  
exp(s(N),X,Y) :- exp(N,X,Z), prod(Z,X,Y).
```