

### Problem 1

Here is an exercise on type compatibility.

Given the following program fragment in some hypothetical language:

```
type S1 is struct {
    int y;
    int w;
};
type S2 is struct {
    int y;
    int w;
};
type S3 is struct {
    int y;
};
S3 f(S1 p) { ... };
...
S1 a, x;
S2 b;
S3 c;
int d;
...

a = b;      // (1)
x = a;      // (2)
c = f(b);   // (3)
d = f(a);   // (4)
```

a) Under name compatibility ("nominal type checking"), which of the four statements (1) ... (4) are type correct (and which are not).

b) Same question under structural compatibility ("structural type checking").

### Problem 2

We have the following classes:

```
class Food {...}
class Cheese extends Food {...}
```

Assume that we have the following functions:

```
int f(c Cheese) {...}

int f'(f Food) {...}
```

someFood is a value of type Food, and someCheese is a value of type Cheese. Then we know that

`f'(someCheese)` can be substituted for `f(someCheese)`

that is, whenever we have a call '`f(someCheese)`' we may just as well call `f'` with the same `someCheese` parameter without causing any static type errors: `f'` can be said to be a subtype of `f`.

Why cannot `f(someFood)` be substituted for `f'(someFood)`? That is why can not `f` be said to be a subtype of `f'`? Give an example of class `Cheese` (that is a more elaborate `Cheese` than above) and a definition of `f` that will create a type error.

### **Problem 3 - Exercise 10.2 in Mitchell's book.**

### **Problem 4**

Consider the classes `C` and `SC` from the lecture slides.

We know that this language allows overloaded methods to be inherited, that is the scope for overloaded methods for a subclass includes the inherited methods.

Here is the answer to the question posed at the lecture (to which method are the different calls bound):

```
C c      = new C();
SC sc    = new SC();
C c'     = new SC();

c.equals(c)    //1    equals 1
c.equals(c')  //2    equals 1
c.equals(sc)  //3    equals 1

c'.equals(c)   //4    equals 1
c'.equals(c') //5    equals 1
c'.equals(sc) //6    equals 1

sc.equals(c)   //7    equals 1
sc.equals(c') //8    equals 1
sc.equals(sc) //9    equals 2
```

It is only in //9 that the `equals 2` method is called, the reason being that overloading (in Java and similar languages) is resolved at compile time. The three calls to `c'` (even though the value of `c'` is a `SC`-object) will be calls to `equals 1`. //7 is also a call to `equals 1`, as the parameter `c` is of type `C` - same with //8.

The method `equals 1` comes in two versions: the `C_equals 1` and the redefined `SC_equals 1`.

a) Indicate for the above first 8 cases which of the `equals 1` are called.

b) Now, suppose that class SC does not have the first equals method, the one with parameter of type C overriding the equals from class C. Determine which of the remaining methods is executed for each of these 8 cases:

```
c.equals(c)      //1
c.equals(c')     //2
c.equals(sc)     //3

c'.equals(c)     //4
c'.equals(c')    //5
c'.equals(sc)    //6

sc.equals(c)     //7
sc.equals(c')    //8
```

### **Problem 5**

a) Write in Java both an abstract data type and a class for the data type Date, with year, month and day, operations before and after and daysBetween. In the abstract data type the operations before, after and daysBetween shall take two Dates, while the operations for the class Date shall have just one Date parameter.

b) There is on 'obvious' way of doing this, where Date is depending on how year, month and day is represented (e.g. as int variables). How would you make Date independent of this representation?