— INF4820 — Algorithms for AI and NLP

Common Lisp Essentials

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Lisp



- ► Conceived in the late 1950s by John McCarthy one of the founding fathers of Al.
- Originally intended as a mathematical formalism.
- A family of high-level languages.
- Several dialects, e.g. Scheme, Clojure, Emacs Lisp, and Common Lisp.
- Although a multi-paradigm language, functional style prevalent.





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- ► the read—eval—print loop.
- ► (= the interactive Lisp-environment)
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- → "this is a string"



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- ? "this is a string"
- \rightarrow "this is a string"
- ? 42
- \rightarrow 42



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- ? "this is a string"
- \rightarrow "this is a string"
- ? 42
- \rightarrow 42
- ? t
- \rightarrow t



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- ? "this is a string"
- \rightarrow "this is a string"
- ? 42
- \rightarrow 42
- ? t
- \rightarrow t
- ? nil \rightarrow nil



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- Symbols evaluate to whatever value they are bound to.

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- \rightarrow t
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- \rightarrow nil
- ? pi
- ightarrow 3.141592653589793d0



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- ? "this is a string"
- \rightarrow "this is a string"
- ? 42
- \rightarrow 42
- ? t
- \rightarrow t
- ? nil
- \rightarrow nil
- ? pi
- \rightarrow 3.141592653589793d0
- ? foo
- → error; unbound

A note on terminology



- ► Lisp manipulates so-called *symbolic expressions*.
- ► AKA s-expressions or sexps.
- ► Two fundamental types of sexps;
 - 1. atoms (e.g., nil, t, numbers, strings, symbols)
 - 2. lists containing other sexps.
- ► Sexps are used to represent *both* data and code.



- "Parenthesized prefix notation"
- First element (prefix) = operator (i.e. the procedure or function).
- The rest of the list is the operands (i.e. the arguments or parameters).
- Use nesting (of lists) to build compound expressions.
- Expressions can span multiple lines; indentation for readability.

- ? (+ 1 2)
- \rightarrow 3



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- ? (+ 1 2)
- \rightarrow 3
- ? (+ 1 2 10 7 5)
- \rightarrow 25



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- ? (+ 1 2 10 7 5)
- \rightarrow 25
- ? (/ (+ 10 20) 2)
- \rightarrow 15



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- ? (+ 1 2)
- \rightarrow 3
- ? (+ 1 2 10 7 5)
- $\rightarrow 25$
- ? (/ (+ 10 20) 2)
- \rightarrow 15
- ? (* (+ 42 58) (- (/ 8 2) 2))
- → 200

The syntax and semantics of CL



- ? (expt (- 8 4) 2)
- \rightarrow 16
 - ▶ You now know (almost) all there is to know about the rules of CL.
 - ► The first element of a list names a function that is invoked with the values of all remaining elements as its arguments.
 - ► A few exceptions, called special forms, with their own evaluation rules.

Creating our own functions



► The special form defun associates a function definition with a symbol:

General form

 $(defun name (parameter_1 ... parameter_n) body)$

Creating our own functions



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General form

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```

```
? (defun average (x y) (/ (+ x y) 2))
```

Creating our own functions



► The special form defun associates a function definition with a symbol:

General form

 $(defun name (parameter_1 ... parameter_n) body)$

- ? (defun average (x y) (/ (+ x y) 2))
- ? (average 10 20)
- \rightarrow 15

Some other special forms



▶ defparameter declares a 'global variable' and assigns a value:

```
? (defparameter *foo* 42) 
? *foo* \rightarrow 42
```

Conditional evaluation with if and cond:

Examples

```
? (if (numberp *foo*)
    "number"
    "something else")
```

→ "number"

General form

```
 \begin{array}{c} (\text{if } \langle predicate \rangle \\ \langle then \ clause \rangle \\ \langle else \ clause \rangle) \end{array}
```

Some other special forms



► defparameter declares a 'global variable' and assigns a value:

```
? (defparameter *foo* 42) 
? *foo* \rightarrow 42
```

► Conditional evaluation with if and cond:

```
Examples
```

```
General form
```

```
(if ⟨predicate⟩
    ⟨then clause⟩
    ⟨else clause⟩)

(cond (⟨predicate₁⟩ ⟨clause₁⟩)
    (⟨predicate₂⟩ ⟨clause₂⟩)
    (⟨predicate₁⟩ ⟨clause₁⟩)
    (t ⟨default clause⟩))
```



$$n! = \begin{cases} 1 & \text{if } n = 0 \\ n \times (n-1)! & \text{if } n > 0 \end{cases}$$



► A recursive procedure: calls itself, directly or indirectly.

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- ► Classic example: the factorial function.
- ► A recursive procedure: calls itself, directly or indirectly.

```
n! = \begin{cases} 1 & \text{if } n = 0\\ n \times (n-1)! & \text{if } n > 0 \end{cases}
```

```
(defun fac (n)
(if (= n 0)
1
(* n (fac (- n 1)))))
```



- Classic example: the factorial function.
- ► A recursive procedure: calls itself, directly or indirectly.
- May seem circular, but is well-defined as long as there's a base case terminating the recursion.

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- Classic example: the factorial function.
- ► A recursive procedure: calls itself, directly or indirectly.
- ► May seem circular, but is well-defined as long as there's a base case terminating the recursion.
- ► For comparison: a non-recursive implementation (in Python).

```
n! = \begin{cases} 1 & \text{if } n = 0 \\ n \times (n-1)! & \text{if } n > 0 \end{cases}
```

(defun fac (n) (if (= n 0) 1 (* n (fac (- n 1)))))

```
def fac(n):
    r = 1
    while (n > 0):
        r = r * n
        n = n - 1
    return r
```



- ► A more efficient way to define *n*! recursively.
- Use a helper procedure with an accumulator variable to collect the product along the way.



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- ▶ no work remains to be done in the calling function.
- ► Once we reach the base case, the return value is ready.



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- Most CL compilers do tail call optimization, so that the recursion is executed as an iterative loop.



- ► A more efficient way to define n! recursively.
- Use a helper procedure with an accumulator variable to collect the product along the way.
- The recursive call is in tail position:

- ▶ no work remains to be done in the calling function.
- ▶ Once we reach the base case, the return value is ready.
- ▶ Most CL compilers do *tail call optimization*, so that the recursion is executed as an iterative loop.
- ► (The next lecture will cover CL's built-in loop construct.)

Tracing the processes



Recursive

```
(defun fac (n)
  (if (= n 0)
        (* n (fac (- n 1)))))
? (fac 7)
\Rightarrow (* 7 (fac 6))
\Rightarrow (* 7 (* 6 (fac 5)))
\Rightarrow (* 7 (* 6 (* 5 (fac 4))))
\Rightarrow (* 7 (* 6 (* 5 (* 4 (fac 3)))))
\Rightarrow (* 7 (* 6 (* 5 (* 4 (* 3 (fac 2))))))
\Rightarrow (* 7 (* 6 (* 5 (* 4 (* 3 (* 2 (fac 1)))))))
\Rightarrow (* 7 (* 6 (* 5 (* 4 (* 3 (* 2 1))))))
\Rightarrow (* 7 (* 6 (* 5 (* 4 (* 3 2)))))
\Rightarrow (* 7 (* 6 (* 5 (* 4 6))))
\Rightarrow (* 7 (* 6 (* 5 24)))
\Rightarrow (* 7 (* 6 120))
\Rightarrow (* 7 720)
\rightarrow 5040
```

Iterative (tail recursive)

```
(defun fac (n)
  (fac-iter 1 1 n))
(defun fac-iter (prod count n)
  (if (> count n)
       prod
       (fac-iter (* count prod)
                    (+ count 1)
                    n)))
? (fac 7)
\Rightarrow (fac-iter 1 1 7)
\Rightarrow (fac-iter 1 2 7)
\Rightarrow (fac-iter 2 3 7)
\Rightarrow (fac-iter 6 4 7)
\Rightarrow (fac-iter 24 5 7)
⇒ (fac-iter 120 6 7)
⇒ (fac-iter 720 7 7)
\Rightarrow (fac-iter 5040 8 7)
\rightarrow 5040
```

Tracing the processes



Recursive

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(defun fac (n)
  (if (= n 0)
        (* n (fac (- n 1)))))
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\Rightarrow (* 7 (* 6 (* 5 24)))
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\Rightarrow (fac-iter 24 5 7)
\Rightarrow (fac-iter 120 6 7)
\Rightarrow (fac-iter 720 7 7)
\Rightarrow (fac-iter 5040 8 7)
\rightarrow 5040
```

The quote operator



- ► A *special form* making expressions self-evaluating.
- ► The quote operator (or simply ''') suppresses evaluation.



- $\,\blacktriangleright\,$ A special form making expressions self-evaluating.
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? $pi \rightarrow 3.141592653589793d0$



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- ► The quote operator (or simply ''') suppresses evaluation.

```
? pi \rightarrow 3.141592653589793d0
```

```
? (quote pi) \rightarrow pi
```



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- ► The quote operator (or simply ''') suppresses evaluation.

```
? pi→ 3.141592653589793d0
```

? (quote pi)
$$\rightarrow$$
 pi

? 'pi
$$\rightarrow$$
 pi



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```
? pi→ 3.141592653589793d0
```

```
? (quote pi) \rightarrow pi
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? 'pi
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 pi

? foobar \rightarrow error; unbound variable



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- ? pi→ 3.141592653589793d0
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- ? foobar \rightarrow error; unbound variable
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```
? pi→ 3.141592653589793d0
```

? (quote pi) \rightarrow pi

? 'pi $\rightarrow pi$

? foobar \rightarrow error; unbound variable

? 'foobar \rightarrow foobar

? (* 2 pi) \rightarrow 6.283185307179586d0



- ► A *special form* making expressions self-evaluating.
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```
? pi→ 3.141592653589793d0
```

- ? (quote pi) \rightarrow pi
- ? 'pi \rightarrow pi
- ? foobar \rightarrow error; unbound variable
- ? 'foobar \rightarrow foobar
- ? (* 2 pi) \rightarrow 6.283185307179586d0
- ? '(* 2 pi) \rightarrow



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```
? pi→ 3.141592653589793d0
```

```
? (quote pi) \rightarrow pi
```

? 'pi
$$\rightarrow$$
 pi

? foobar
$$\rightarrow$$
 error; unbound variable

? (* 2 pi)
$$\rightarrow$$
 6.283185307179586d0

?'(* 2 pi)
$$\rightarrow$$
 (* 2 pi)



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? pi \rightarrow 3.141592653589793d0
```

```
? (quote pi) \rightarrow pi
```

$$?$$
 'pi \rightarrow pi

? foobar
$$\rightarrow$$
 error; unbound variable

? (* 2 pi)
$$\rightarrow$$
 6.283185307179586d0

?'(* 2 pi)
$$\rightarrow$$
 (* 2 pi)

?()
$$\rightarrow$$
 error; missing procedure



- ► A special form making expressions self-evaluating.
- ► The quote operator (or simply ''') suppresses evaluation.
- ? $pi \rightarrow 3.141592653589793d0$
- ? (quote pi) \rightarrow pi
- ? 'pi \rightarrow pi
- ? foobar \rightarrow error; unbound variable
- ? 'foobar \rightarrow foobar
- ? '(* 2 pi) \rightarrow (* 2 pi)
- ? () \rightarrow error; missing procedure

? (* 2 pi) \rightarrow 6.283185307179586d0

? ¹() →



- ► A special form making expressions self-evaluating.
- ► The quote operator (or simply ''') suppresses evaluation.
- ? $pi \rightarrow 3.141592653589793d0$
- ? (quote pi) \rightarrow pi
- ? 'pi → pi
- ? foobar \rightarrow error; unbound variable
- ? 'foobar → foobar
- . 100bai / 100bai
- ?'(* 2 pi) \rightarrow (* 2 pi)
- ? () \rightarrow error; missing procedure

? (* 2 pi) \rightarrow 6.283185307179586d0

? ¹() → ()

Both code and data are s-expressions



- ▶ We've mentioned how sexps are used to represent *both* data and code.
- ► Note the double role of lists:
- ► Lists are function calls:

```
? (* 10 (+ 2 3)) → 50
```

? (bar 1 2) \rightarrow error; function bar undefined

Both code and data are s-expressions



- ▶ We've mentioned how sexps are used to represent *both* data and code.
- ► Note the double role of lists:
- ► Lists are function calls;

```
? (* 10 (+ 2 3)) \rightarrow 50
```

? (bar 1 2) \rightarrow error; function bar undefined

▶ But, lists can also be data;

```
? '(foo bar) \rightarrow (foo bar)
? (list 'foo 'bar) \rightarrow (foo bar)
```

Break









http://xkcd.com/297/

Eric Raymond, How to Become a Hacker, 2001:

Lisp is worth learning for the profound enlightenment experience you will have when you finally get it;

that experience will make you a better programmer for the rest of your days, even if you should never actually use Lisp itself a lot.



```
? (cons 1 (cons 2 (cons 3 nil))) \rightarrow (1 2 3)
```



```
? (cons 1 (cons 2 (cons 3 nil))) \rightarrow (1 2 3)
? (cons 0 '(1 2 3)) \rightarrow
```



```
? (cons 1 (cons 2 (cons 3 nil))) \rightarrow (1 2 3)
? (cons 0 '(1 2 3)) \rightarrow (0 1 2 3)
```



```
? (cons 1 (cons 2 (cons 3 nil))) \rightarrow (1 2 3)
? (cons 0 '(1 2 3)) \rightarrow (0 1 2 3)
? (first '(1 2 3)) \rightarrow 1
```



```
? (cons 1 (cons 2 (cons 3 nil))) \rightarrow (1 2 3)
? (cons 0 '(1 2 3)) \rightarrow (0 1 2 3)
? (first '(1 2 3)) \rightarrow 1
? (rest '(1 2 3)) \rightarrow (2 3)
```



```
? (cons 1 (cons 2 (cons 3 nil))) \rightarrow (1 2 3)

? (cons 0 '(1 2 3)) \rightarrow (0 1 2 3)

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? (rest '(1 2 3)) \rightarrow (2 3)

? (first (rest '(1 2 3))) \rightarrow
```



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? (cons 1 (cons 2 (cons 3 nil))) \rightarrow (1 2 3)

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? (rest '(1 2 3)) \rightarrow (2 3)

? (first (rest '(1 2 3))) \rightarrow 2
```



```
? (cons 1 (cons 2 (cons 3 nil))) → (1 2 3)
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? (rest '(1 2 3)) → (2 3)
? (first (rest '(1 2 3))) → 2
? (rest (rest (rest '(1 2 3)))) →
```



```
? (cons 1 (cons 2 (cons 3 nil))) → (1 2 3)
? (cons 0 '(1 2 3)) → (0 1 2 3)
? (first '(1 2 3)) → 1
? (rest '(1 2 3)) → (2 3)
? (first (rest '(1 2 3))) → 2
? (rest (rest (rest '(1 2 3)))) → nil
```



► cons builds up new lists; first and rest destructure them.

```
? (cons 1 (cons 2 (cons 3 nil))) \rightarrow (1 2 3)

? (cons 0 '(1 2 3)) \rightarrow (0 1 2 3)

? (first '(1 2 3)) \rightarrow 1

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? (first (rest '(1 2 3))) \rightarrow 2

? (rest (rest (rest '(1 2 3)))) \rightarrow nil
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```
? (list 1 2 3) \rightarrow (1 2 3)
```



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? (rest (rest (rest '(1 2 3)))) \rightarrow nil
```

```
? (list 1 2 3) \rightarrow (1 2 3)
? (append '(1 2) '(3) '(4 5 6)) \rightarrow (1 2 3 4 5 6)
```



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```

```
? (list 1 2 3) \rightarrow (1 2 3) 
? (append '(1 2) '(3) '(4 5 6)) \rightarrow (1 2 3 4 5 6) 
? (length '(1 2 3)) \rightarrow 3
```



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? (cons 1 (cons 2 (cons 3 nil))) \rightarrow (1 2 3)

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? (first (rest '(1 2 3))) \rightarrow 2

? (rest (rest (rest '(1 2 3)))) \rightarrow nil
```

```
? (list 1 2 3) → (1 2 3)
? (append '(1 2) '(3) '(4 5 6)) → (1 2 3 4 5 6)
? (length '(1 2 3)) → 3
? (reverse '(1 2 3)) → (3 2 1)
```



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? (cons 1 (cons 2 (cons 3 nil))) \rightarrow (1 2 3)

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? (list 1 2 3) → (1 2 3)
? (append '(1 2) '(3) '(4 5 6)) → (1 2 3 4 5 6)
? (length '(1 2 3)) → 3
? (reverse '(1 2 3)) → (3 2 1)
? (nth 2 '(1 2 3)) → 3
```



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? (reverse '(1 2 3)) → (3 2 1)
? (nth 2 '(1 2 3)) → 3
? (last '(1 2 3)) → (3)
```



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? (cons 1 (cons 2 (cons 3 nil))) \rightarrow (1 2 3)

? (cons 0 '(1 2 3)) \rightarrow (0 1 2 3)

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? (first (rest '(1 2 3))) \rightarrow 2

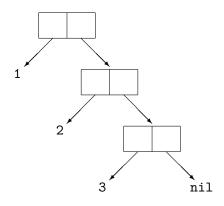
? (rest (rest (rest '(1 2 3)))) \rightarrow nil
```

```
? (list 1 2 3) → (1 2 3)
? (append '(1 2) '(3) '(4 5 6)) → (1 2 3 4 5 6)
? (length '(1 2 3)) → 3
? (reverse '(1 2 3)) → (3 2 1)
? (nth 2 '(1 2 3)) → 3
? (last '(1 2 3)) → (3) Wait, why not 3?
```

Lists are really chained 'cons cells'



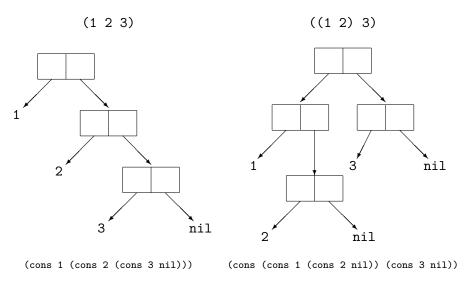




(cons 1 (cons 2 (cons 3 nil)))

Lists are really chained 'cons cells'





Assigning values: 'Generalized variables'



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```



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```



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```



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```



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? (setf *foo* (+ *foo* 1))
? *foo* → 43
? (setf *foo* '(2 2 3))
? (setf (first *foo*) 1)
? *foo* → (1 2 3)
```

Some other macros for assignment



| Example | Type of x | Effect |
|---------------|-----------|--|
| (incf x y) | number | (setf x (+ x y)) |
| (incf x) | number | (incf x 1) |
| (decf x y) | number | (setf x (- x y)) |
| (decf x) | number | (decf x 1) |
| (push y x) | list | (setf x (cons y x)) |
| (pop x) | list | <pre>(let ((y (first x))) (setf x (rest x)) y)</pre> |
| (pushnew y x) | list | <pre>(if (member y x) x (push y x))</pre> |



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```
? (defparameter *foo* 42)
? (defparameter *bar* 100)
? (let ((*bar* 7)
         (baz 1))
    (+ baz *bar* *foo*))
\rightarrow 50
? *bar* \to 100
? baz \rightarrow
```



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         (baz 1))
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\rightarrow 50
? *bar* \to 100
? baz → error; unbound variable
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- ▶ Previously existing bindings are *shadowed* within the lexical scope.



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- ▶ Bindings valid only in the body of let.
- ▶ Previously existing bindings are *shadowed* within the lexical scope.
- ▶ let* is like let but binds sequentially.

Predicates



- ► A *predicate* tests some condition.
- ► Evaluates to a boolean truth value:
 - ▶ nil (the empty list) means false.
 - ► Anything non-nil (including t) means *true*.

▶ Plethora of equality tests: eq, eq1, equal, and equalp.



- eq tests object identity; it is not useful for numbers or characters.
- ▶ eql is like eq, but well-defined on numbers and characters.
- equal tests structural equivalence
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```
? (eq (list 1 2 3) '(1 2 3)) \rightarrow nil
```



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```
? (eq (list 1 2 3) '(1 2 3)) \to nil
? (equal (list 1 2 3) '(1 2 3)) \to t
```



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? (eq (list 1 2 3) '(1 2 3)) \rightarrow nil
? (equal (list 1 2 3) '(1 2 3)) \rightarrow t
? (eq 42 42) \rightarrow ? [implementation-dependent]
```



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? (equal (list 1 2 3) '(1 2 3)) \rightarrow t
? (eq 42 42) \rightarrow ? [implementation-dependent]
? (eq1 42 42) \rightarrow t
```



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```
? (eq (list 1 2 3) '(1 2 3)) \rightarrow nil
? (equal (list 1 2 3) '(1 2 3)) \rightarrow t
? (eq 42 42) \rightarrow ? [implementation-dependent]
? (eql 42 42) \rightarrow t
? (eql 42 42.0) \rightarrow nil
```



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? (eq1 42 42.0) \rightarrow nil
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```



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? (eq (list 1 2 3) '(1 2 3)) \rightarrow nil
? (equal (list 1 2 3) '(1 2 3)) \rightarrow t
? (eq 42 42) \rightarrow ? [implementation-dependent]
? (eq1 42 42) \rightarrow t
? (eq1 42 42.0) \rightarrow nil
? (equalp 42 42.0) \rightarrow t
? (equal "foo" "foo") \rightarrow t
```



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```
? (eq (list 1 2 3) '(1 2 3)) \rightarrow nil

? (equal (list 1 2 3) '(1 2 3)) \rightarrow t

? (eq 42 42) \rightarrow ? [implementation-dependent]

? (eq1 42 42) \rightarrow t

? (eq1 42 42.0) \rightarrow nil

? (equalp 42 42.0) \rightarrow t

? (equal "foo" "foo") \rightarrow t

? (equalp "FOO" "foo") \rightarrow t
```



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? (eq1 42 42) \rightarrow t
? (eq1 42 42.0) \rightarrow nil
? (equalp 42 42.0) \rightarrow t
? (equal "foo" "foo") \rightarrow t
? (equalp "F00" "foo") \rightarrow t
```

► Also many type-specialized tests like =, string=, etc.



- ► Symbols can have values as functions and variables at the same time.
- ▶ #' (sharp-quote) gives us the function object bound to a symbol.

```
? (defun foo (x)
(* x 1000))
```



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```
? (defun foo (x)
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```

? (defparameter foo 42) \rightarrow 2



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```
? (defun foo (x) (* x 1000))
? (defparameter foo 42) \rightarrow 2
? (foo foo) \rightarrow
```



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```
? (defun foo (x)
          (* x 1000))
? (defparameter
```

? (defparameter foo 42) \rightarrow 2

? (foo foo) \rightarrow 42000



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```
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```

? (defparameter foo 42) \rightarrow 2

? (foo foo) \rightarrow 42000

? foo \rightarrow 42



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```
? (defun foo (x)
(* x 1000))
```

? (defparameter foo 42) \rightarrow 2

? (foo foo) \rightarrow 42000

? foo $\rightarrow 42$

? #'foo → #<Interpreted Function F00>



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- ▶ #' (sharp-quote) gives us the function object bound to a symbol.

```
? (defun foo (x)
     (* x 1000))
? (defparameter foo 42) \rightarrow 2
? (foo foo) \rightarrow 42000
? foo \rightarrow 42
? #'foo → #<Interpreted Function FOO>
? (funcall #'foo foo) \rightarrow 42000
```



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```
? (defun foo (x)
          (* x 1000))
```

- ? (defparameter foo 42) \rightarrow 2
- ? (foo foo) \rightarrow 42000
- ? foo $\rightarrow 42$
- ? #'foo \rightarrow #<Interpreted Function FOO>
- ? (funcall #'foo foo) \rightarrow 42000
- #' and funcall (as well as apply) are useful when passing around functions as arguments.



- ► Functions that accept functions as arguments or return values.
- ► Functions in Lisp are first-class objects.
 - ► Can be created at run-time, passed as arguments, returned as values, stored in variables. . . just like any other data type.



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```
? (defun filter (list test)
   (cond ((null list) nil)
         ((funcall test (first list))
           (cons (first list)
                 (filter (rest list) test)))
         (t (filter (rest list) test))))
? (defparameter foo '(11 22 33 44 55))
? (filter foo #'evenp)
\rightarrow (22 44)
```

Anonymous functions



- ► We can also pass function arguments without first binding them to a name, using lambda expressions: (lambda (parameters) body)
- ► A function definition without the defun and *symbol* part.

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Anonymous functions



- ► We can also pass function arguments without first binding them to a name, using lambda expressions: (lambda (parameters) body)
- ► A function definition without the defun and *symbol* part.

- ► Typically used for ad-hoc functions that are only locally relevant and simple enough to be expressed inline.
- ► Or, when constructing functions as return values.

Returning functions



- ► We have seen how to create anonymous functions using lambda and pass them as arguments.
- ► Now let's combine that with a function that itself returns another function (which we then bind to a variable).

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- ► Now let's combine that with a function that itself returns another function (which we then bind to a variable).

```
(defparameter foo '(11 22 33 44 55))
 (defun make-range-test (lower upper)
    #'(lambda (x)
        (and (> x lower)
              (< x upper))))
? (filter foo (make-range-test 10 30))
\rightarrow (11 22)
```

Programming in INF4820



- ► In the IFI Linux environment, we have available Allegro Common Lisp, a commercial Lisp interpreter and compiler.
- ► We will provide a pre-configured, integrated setup with emacs and the SLIME Lisp interaction mode.
- Several open-source Lisp implementations exist, e.g. Clozure or SBCL; compatible with SLIME, so feel free to experiment (at some later point).
- ► First-time users, please spend some time studying basic keyboard commands, for example: C-h t and M-x doctor RET.
- ► See the getting started guide and emacs cheat sheet on the course page.
- ▶ Obligatory assignment 1 is out now, and due Wed. 9th Sept.
 - ► See course page or just run 'svn update'.

Next week



More Common Lisp.

- ► More on argument lists (optional arguments, keywords, defaults).
- ► More data types: Hash-tables, a-lists, arrays, sequences, and structures
- ► More higher-order functions.
- ► Iteration (loop) and mapping.